



**black hat**<sup>®</sup>  
USA 2024

**AUGUST 7-8, 2024**  
BRIEFINGS

# **Bypassing ARM's Memory Tagging Extension with a Side-Channel Attack**

Speaker: Juhee Kim

# Whoami

## Juhee Kim

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Seoul National University

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Focuses on

- Software and Systems security
- Bug finding, Attack mitigation
- Linux kernel, Web browser, GPU/ML systems

# Contributors

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- Security researcher at Samsung Research
- System security, Confidential Computing
- Published in USENIX Security and ASPLOS

## Sihyeon Roh

- Ph.D Student at CompSec Lab
- Hardware side-channels

## Jaeyoung Chung

- Ph.D Student at CompSec Lab
- System Security
- CTF player

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- Ph.D Student at CompSec Lab
- Fuzzing, Browser security, Bug bounty
- CTF player



Samsung Research



Georgia Institute  
of Technology

## Taesoo Kim

- Vice president of Samsung Research
- Professor of Georgia Tech
- Won several best paper awards from USENIX Security, EuroSys

## Byoungyoung Lee

- Professor of Seoul National University
- Leads CompSec Lab
- System security, Confidential computing
- Previous CTF player
- Spoken at Black Hat

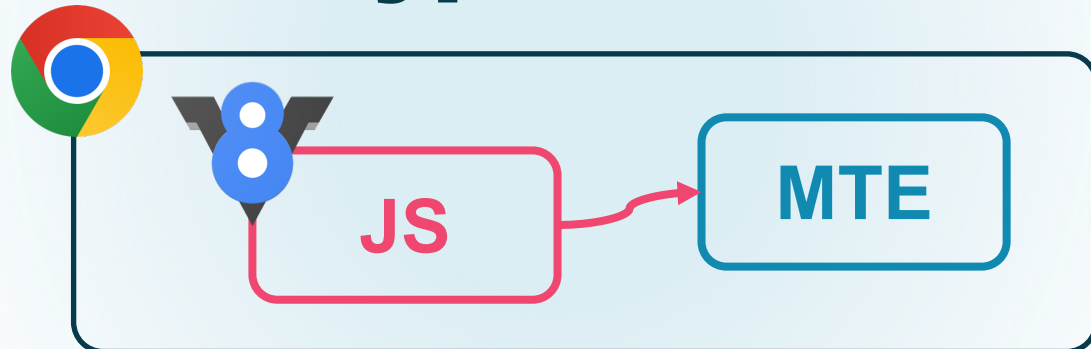
# Roadmap

## ARM Memory Tagging Extension

arm



## Real-world MTE Bypass Attack



## Cache Side-Channel

Cache



## Speculative Execution

if (cond)

True False



## MTE Tag Leakage Side-Channel



MTE



# Roadmap

## ARM Memory Tagging Extension

arm



## Cache Side-Channel

Cache



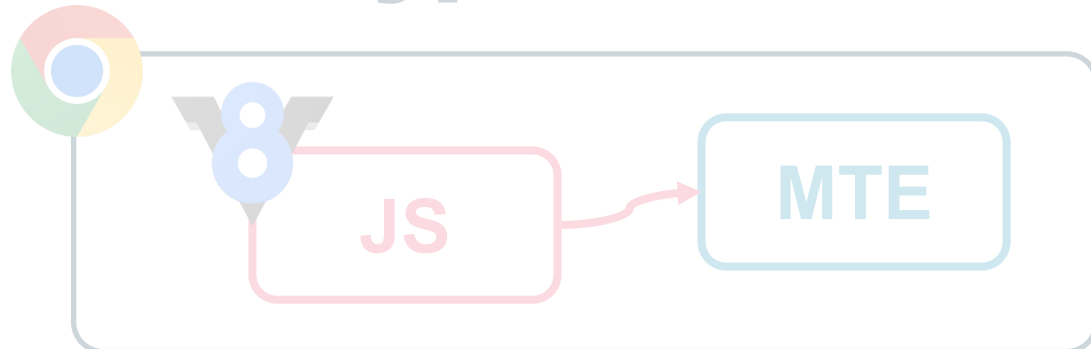
## Speculative Execution

if (cond)

True False



## Real-world MTE Bypass Attack



## MTE Tag Leakage Side-Channel



# Memory corruption attacks

have been the most pervasive and dangerous security threats

## Heartbleed (2014)

OpenSSL information leak

## Bad Binder (2019)

Bad Binder: Android In-The-Wild Exploit

## reggreSSHion (2024)

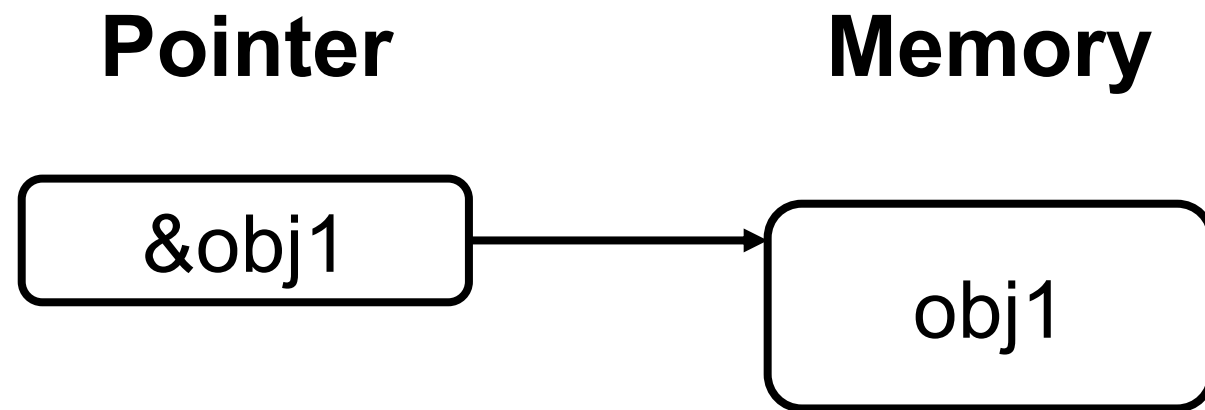
regreSSHion: Remote Unauthenticated Code Execution Vulnerability in OpenSSH server

## BLASTPASS (2023)

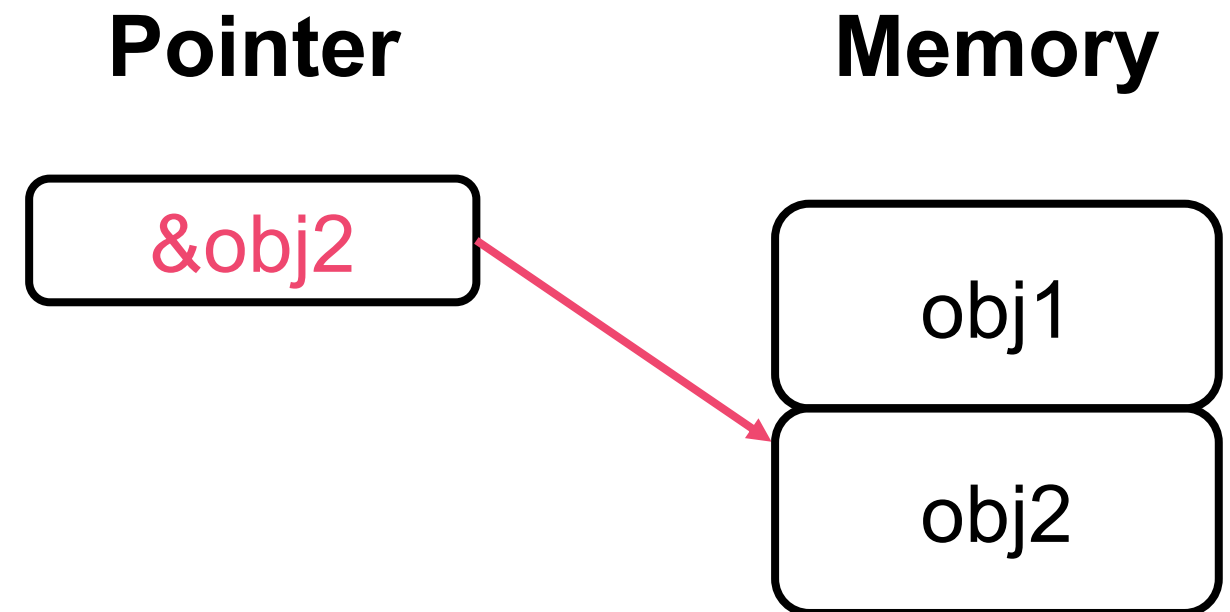
NSO Group iPhone Zero-Click, Zero-Day Exploit Captured in the Wild

# What is Memory Corruption?

## Valid Access



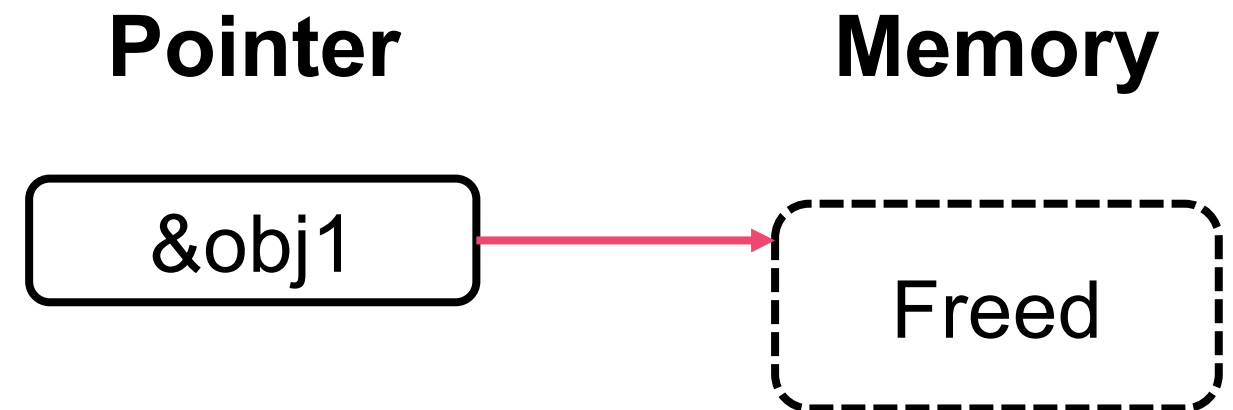
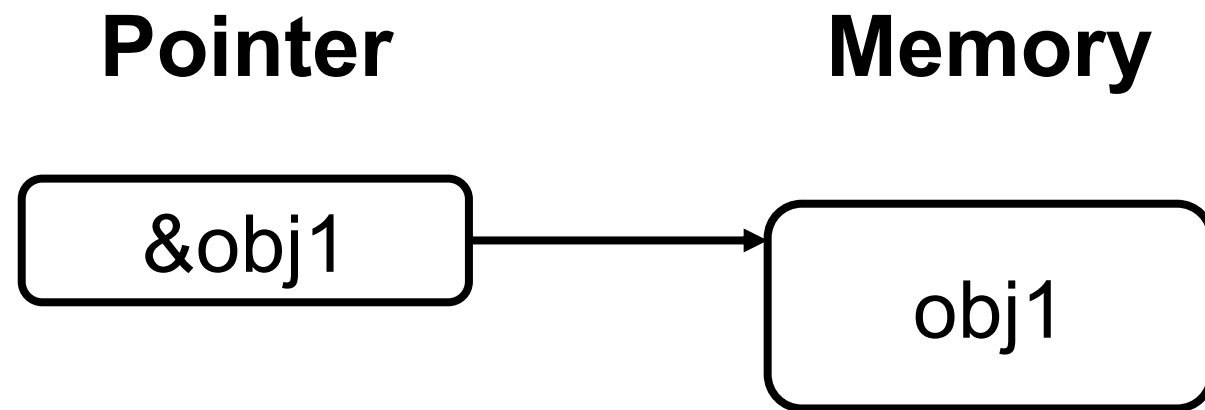
## Invalid Access (Out-of-bounds)



# What is Memory Corruption?

Valid Access

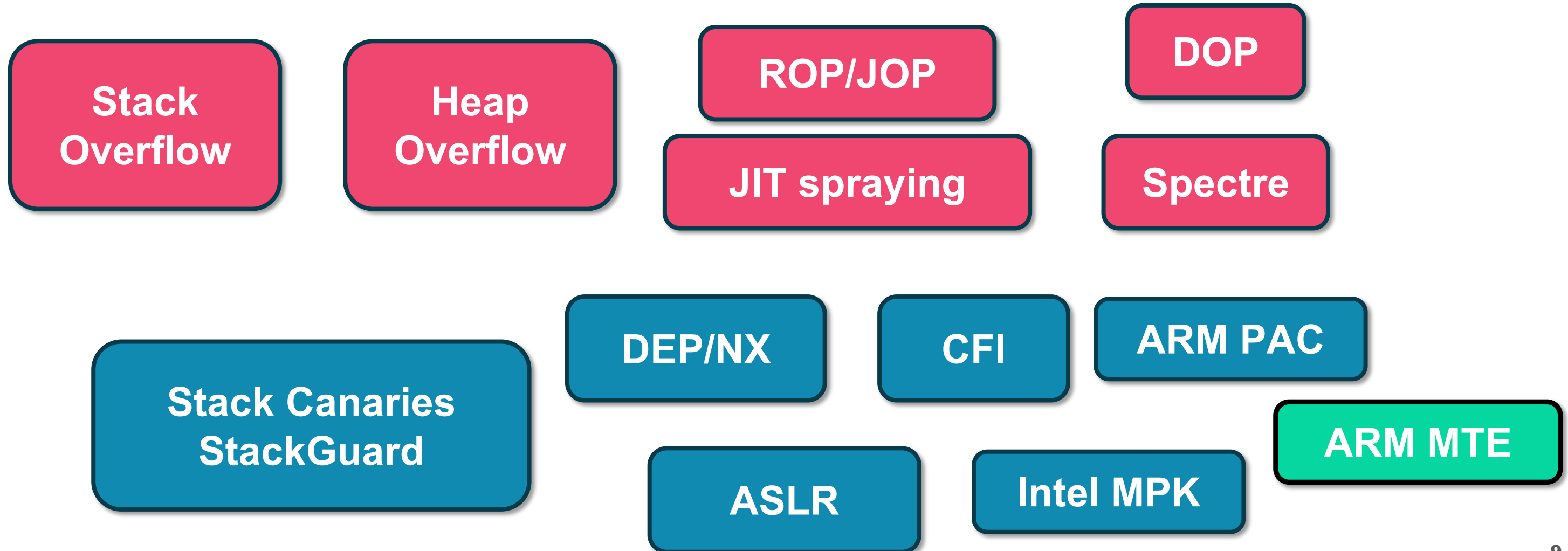
Invalid Access  
(Use-after-free)





# Attack and Defense Techniques

70s-80s — 90s — 2000s — 2010s — 2020s →



# Google Pixel 8 / 8 pro — First MTE hardware released in Sep. 2023



*“MTE being one **key** feature that is delivering **secure mobile experiences**”*

*- Arm (Feb 2023)*

*“**MTE** is still by far the most promising path forward for improving C/C++ software security”*

*- Google Project Zero (Aug 2023)*

*“**Memory tagging** has the potential to provide good value both for **discovering vulnerabilities** and as a **mitigation for vulnerabilities**”*

*- Microsoft (Mar 2020)*

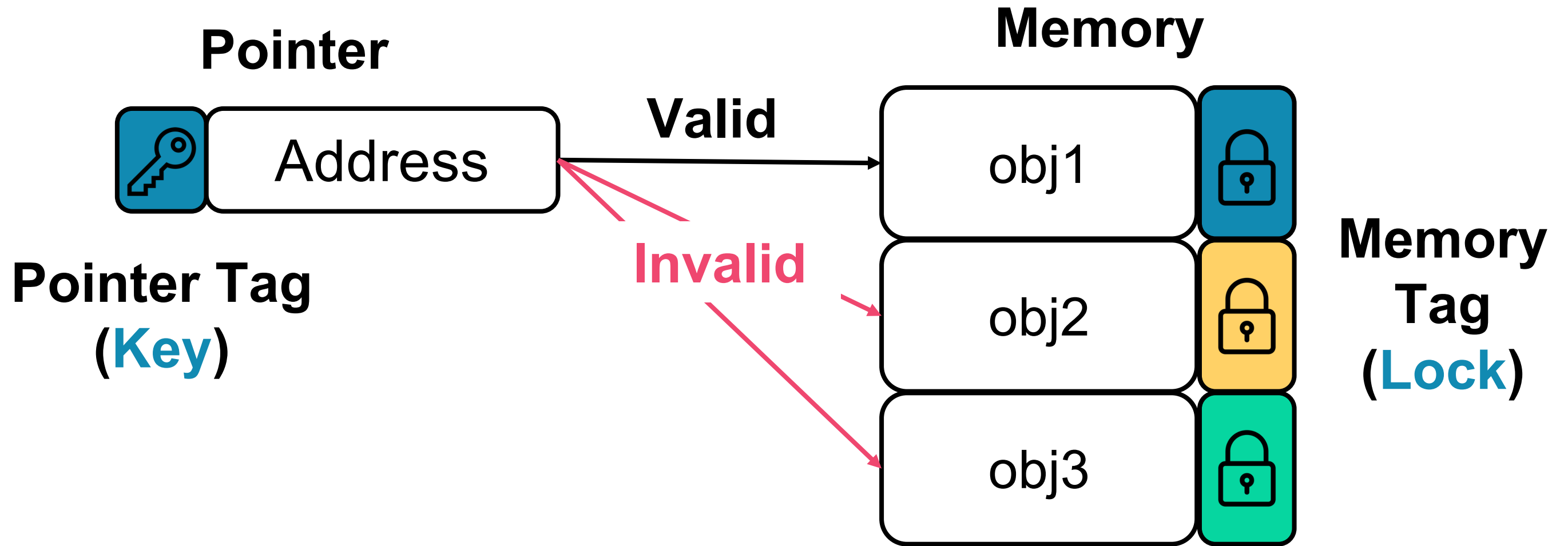
# Why is MTE so Special?

**Hardware-based**

**Memory Corruption Detection**

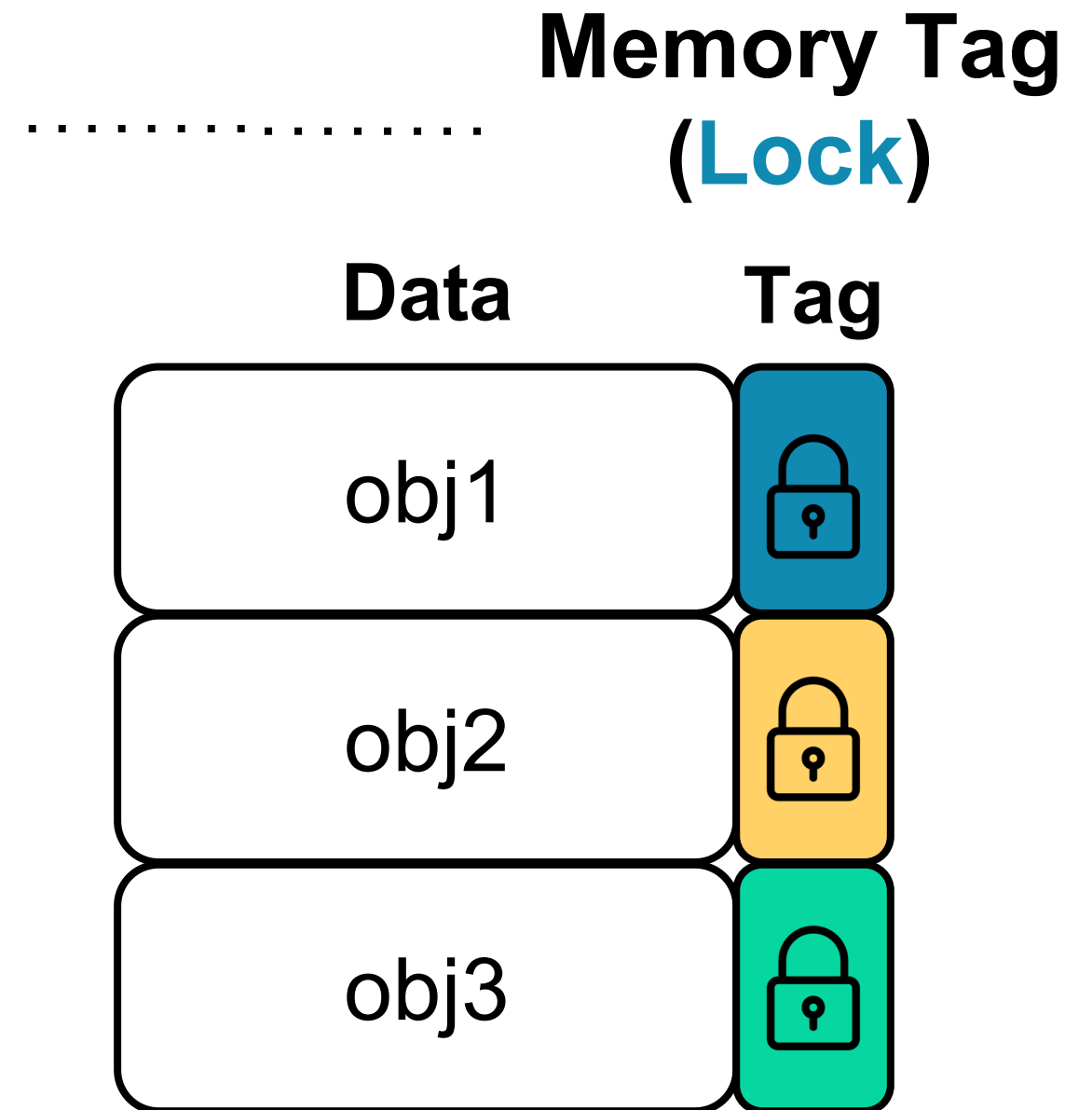
**Fast and Compatible**

# ARM Memory Tagging Extensions



# (1) Memory Tag

Dedicated memory region stores  
a 4-bit tag per 16-byte data



## (2) Pointer Tag

Pointer



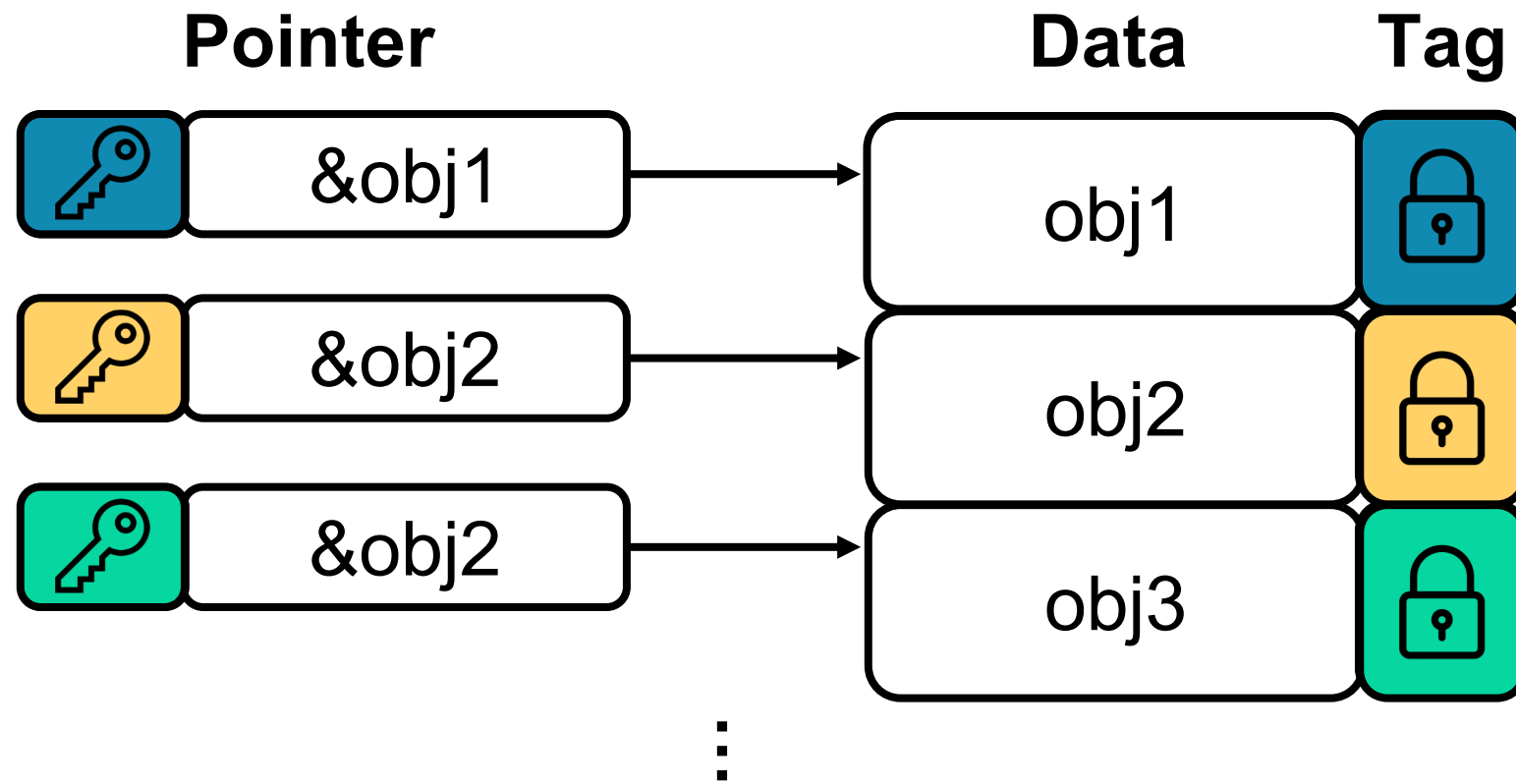
Pointer Tag  
(Key)

.....

**A pointer stores a 4-bit tag in its unused space**

# (3) Tag Allocation

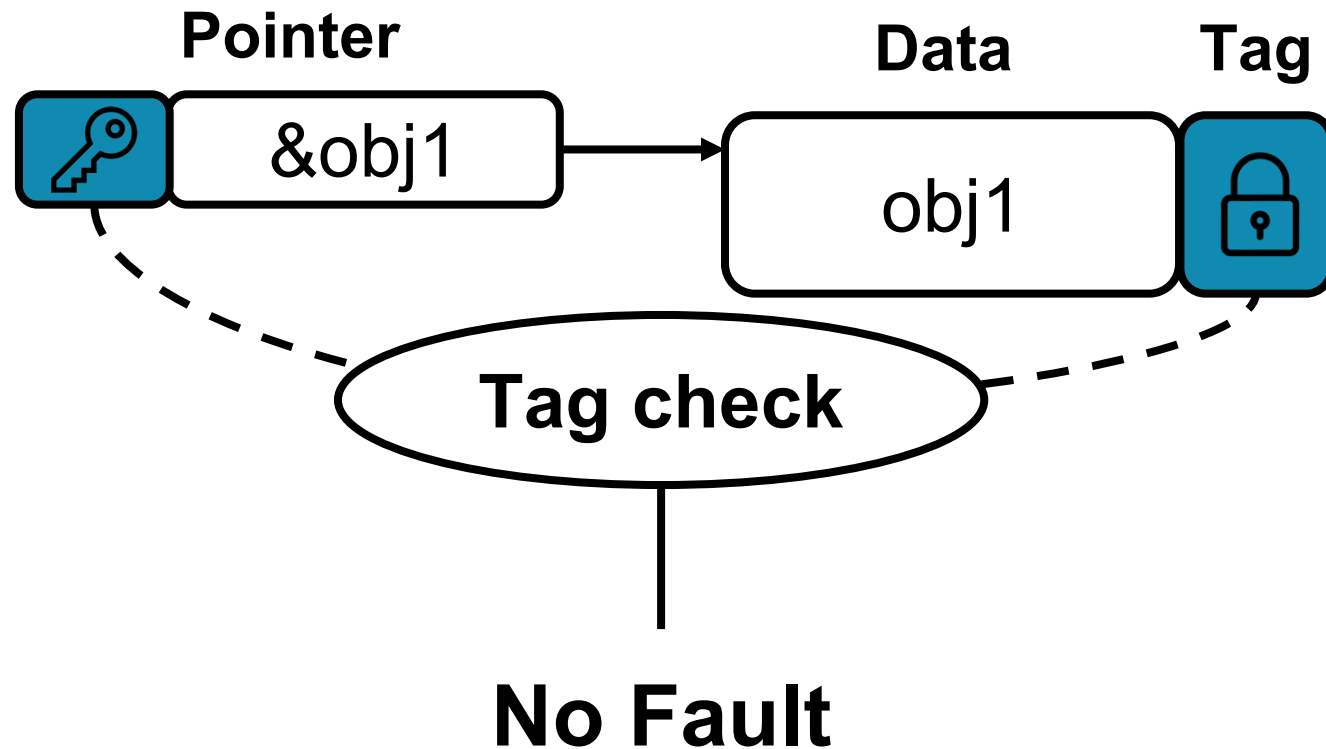
New instructions to **create a random tag** and **load/store memory tags**



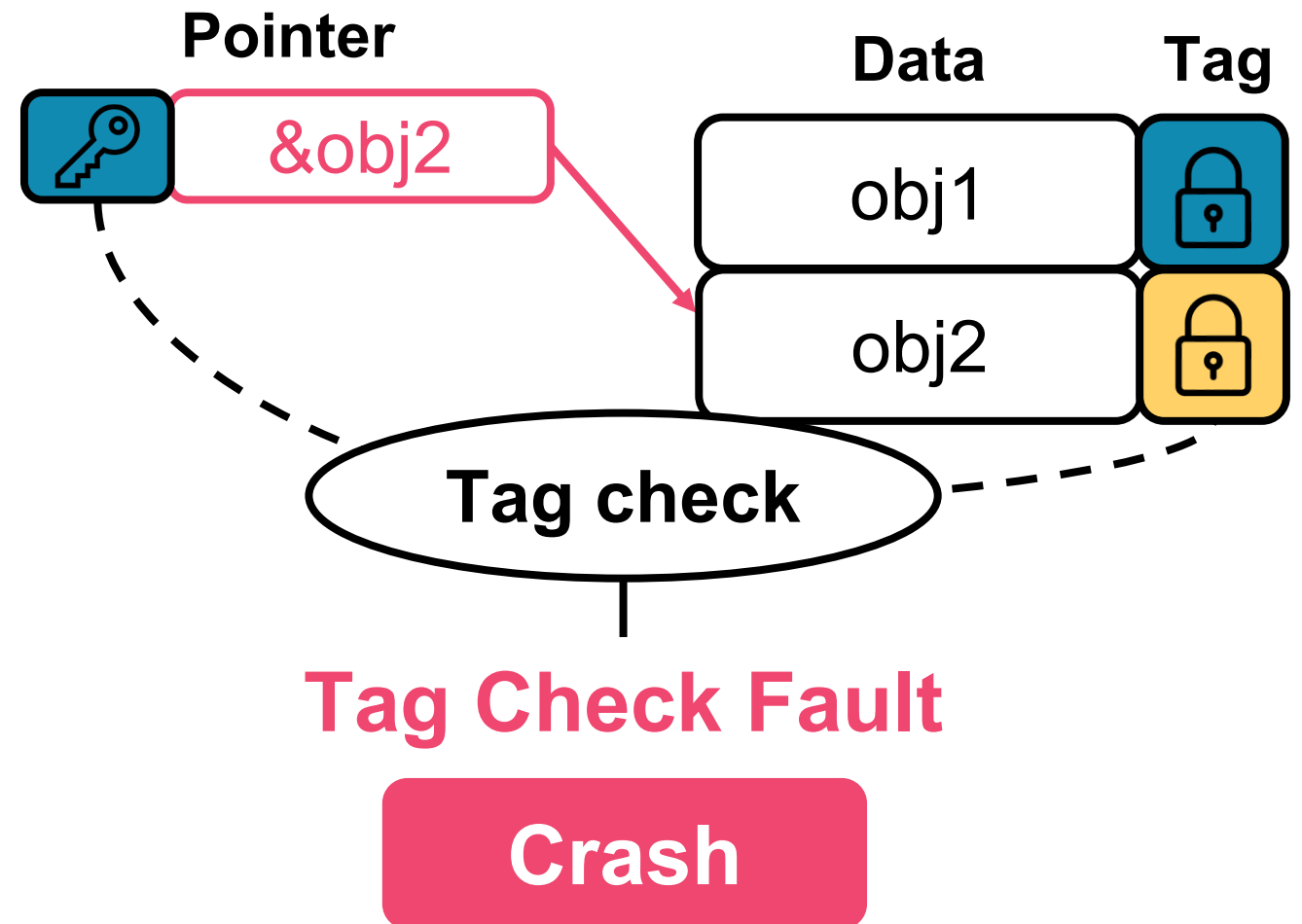
# (4) Tag Check

Transparently done by hardware

**Valid** memory access



**Invalid** memory access

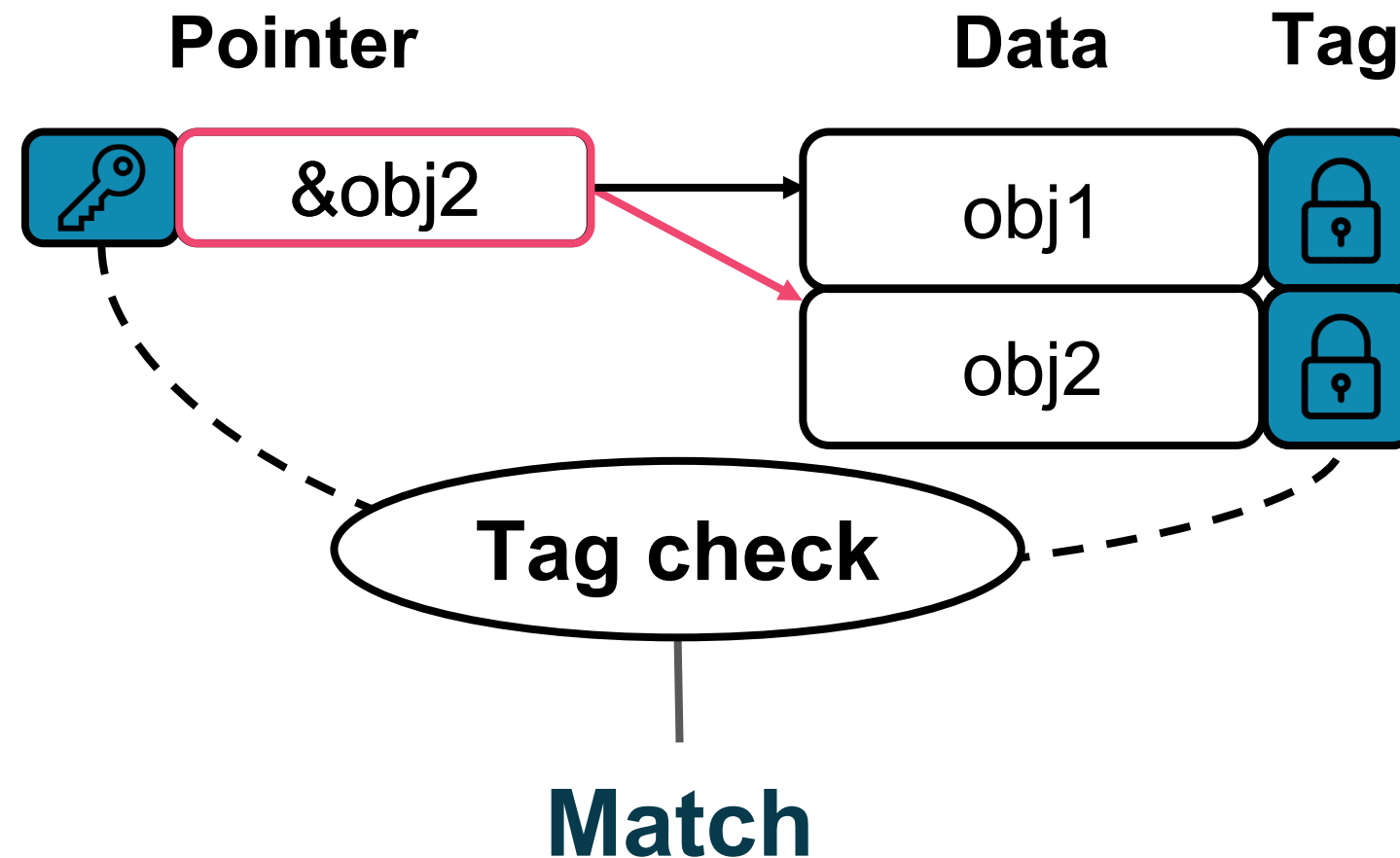




# How to Bypass MTE?

## (1) Tag Collision (16 possible tags)

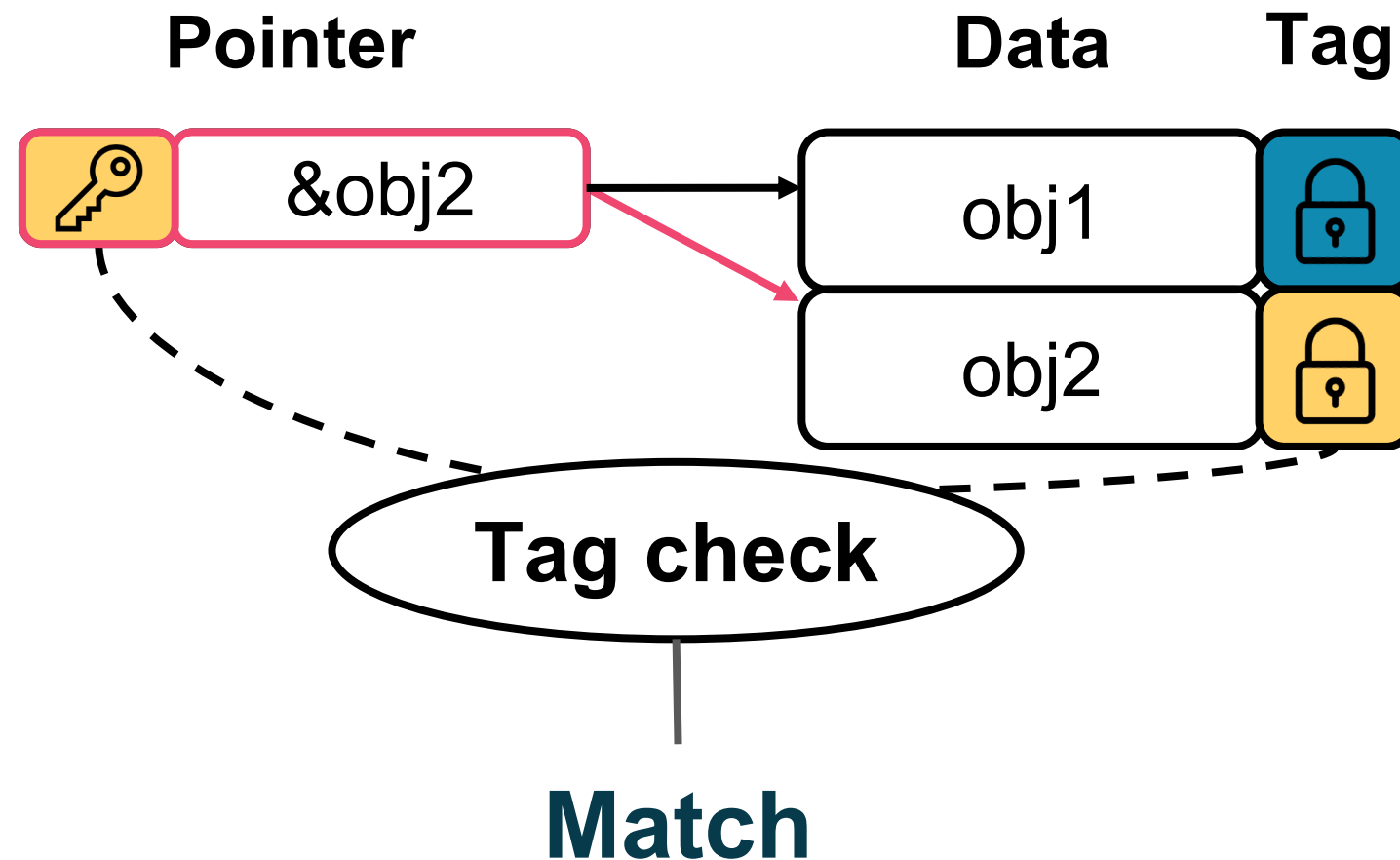
Wait until the **pointer tag** matches the **target memory tag**



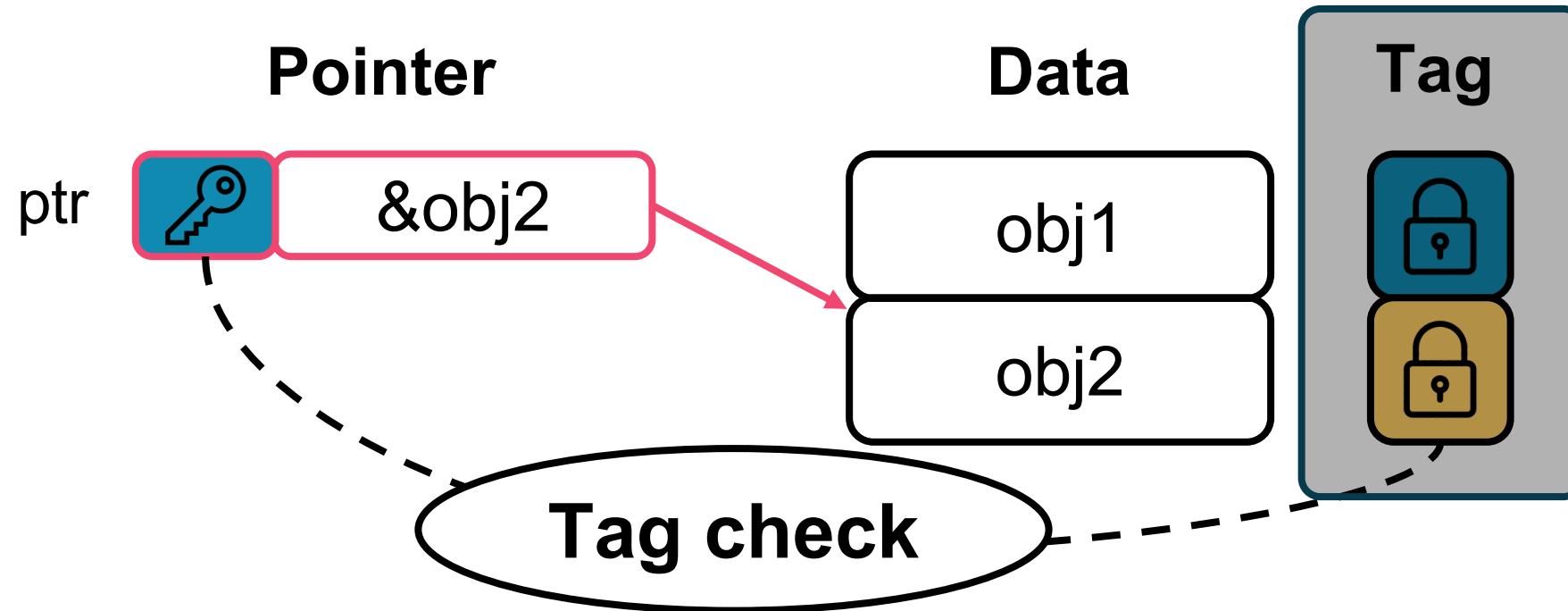
# How to Bypass MTE?

## (2) Pointer Tag Corruption

Corrupt the **pointer tag** to the **target memory tag**



# Challenge: Random Tags



**Match**  $\Rightarrow$  **Attack Succeeds 1/16 (6%)**

**Mismatch**  
**Crash**

$\Rightarrow$  **Attack Fails 15/16 (94%)**

# MTE Bypass Requirement

A Reliable way to **leak MTE tag**  
of any address

# Approach

- Leak tag check result from **Cache Side-channel**
- Exploit **Speculative Execution** to avoid crash

# Roadmap

ARM Memory  
Tagging Extension

arm



Cache  
Side-Channel

Cache



Speculative  
Execution

if (cond)

True False



Real-world  
MTE Bypass Attack



JS

MTE

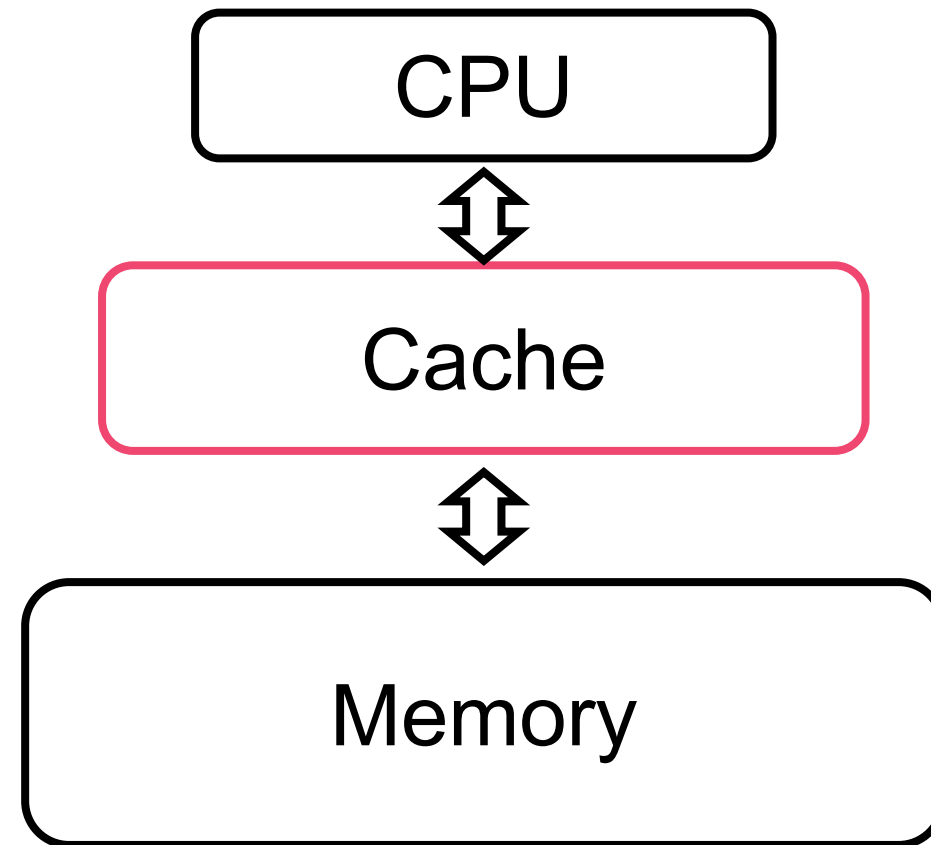
MTE Tag Leakage  
Side-Channel



MTE



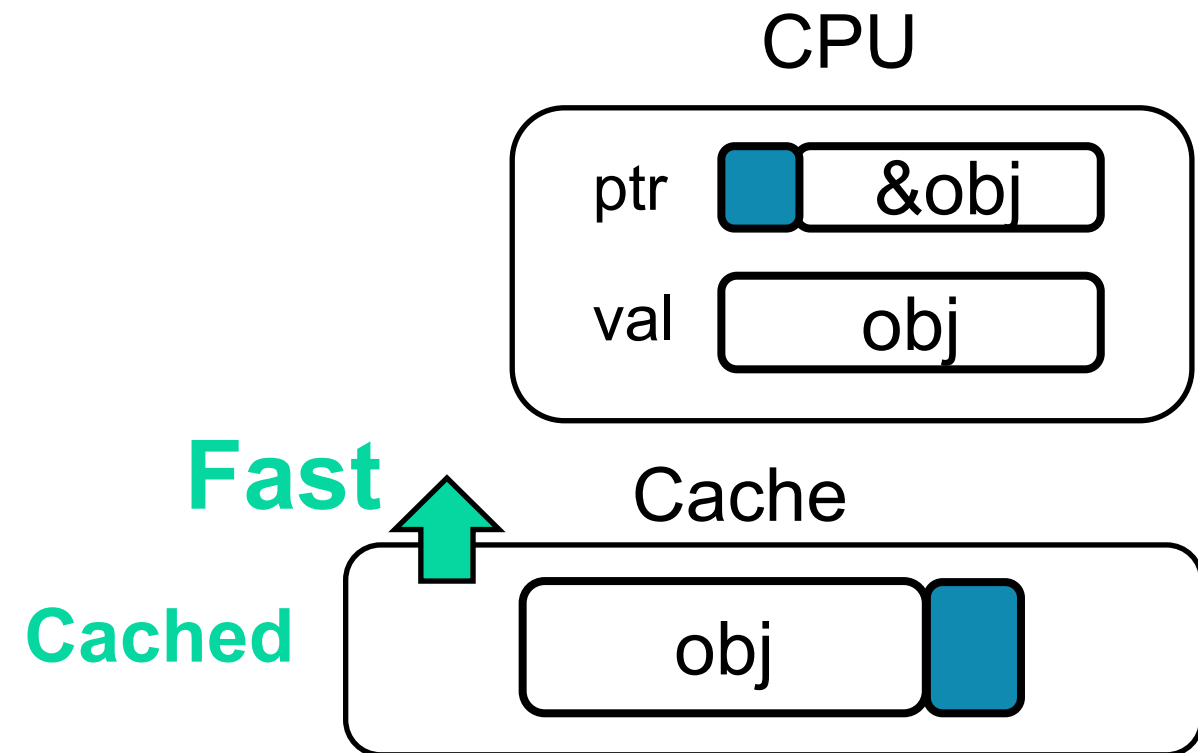
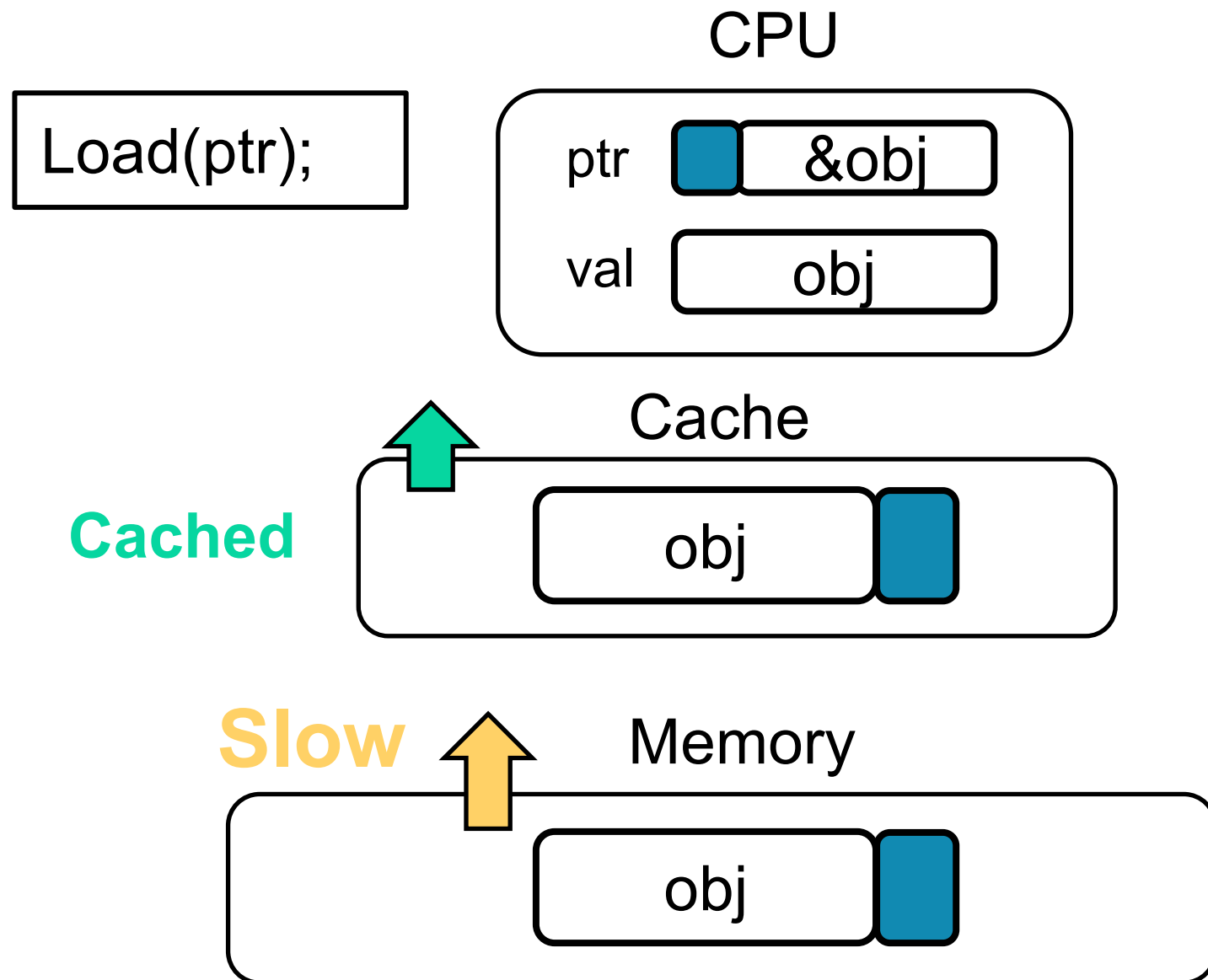
# What is Cache?



# What is Cache?

**First Access : Slow**

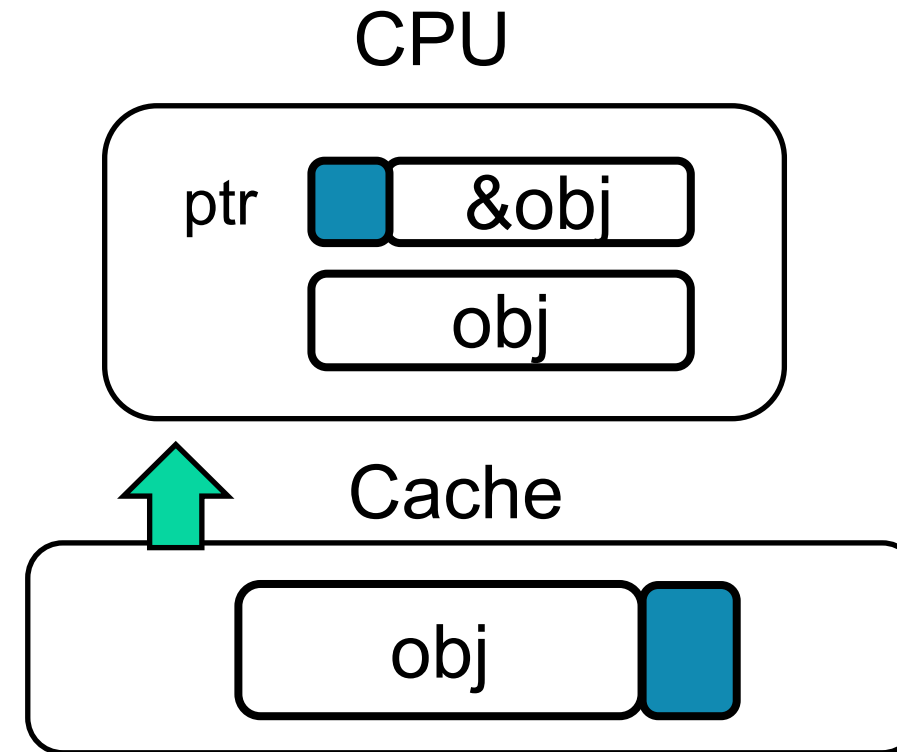
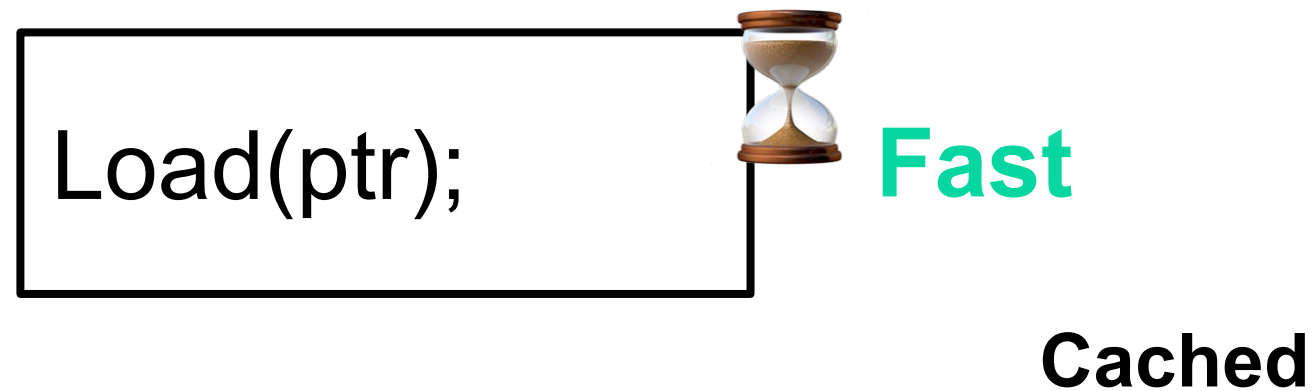
**Second Access : Fast**





# Cache Side-Channel

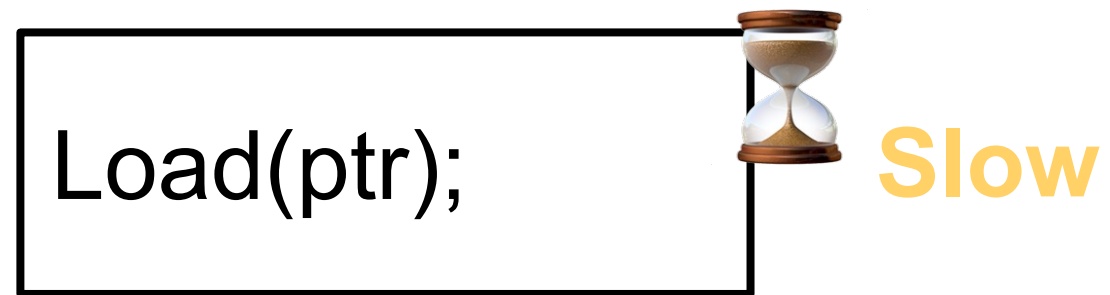
Q. Has obj been accessed?



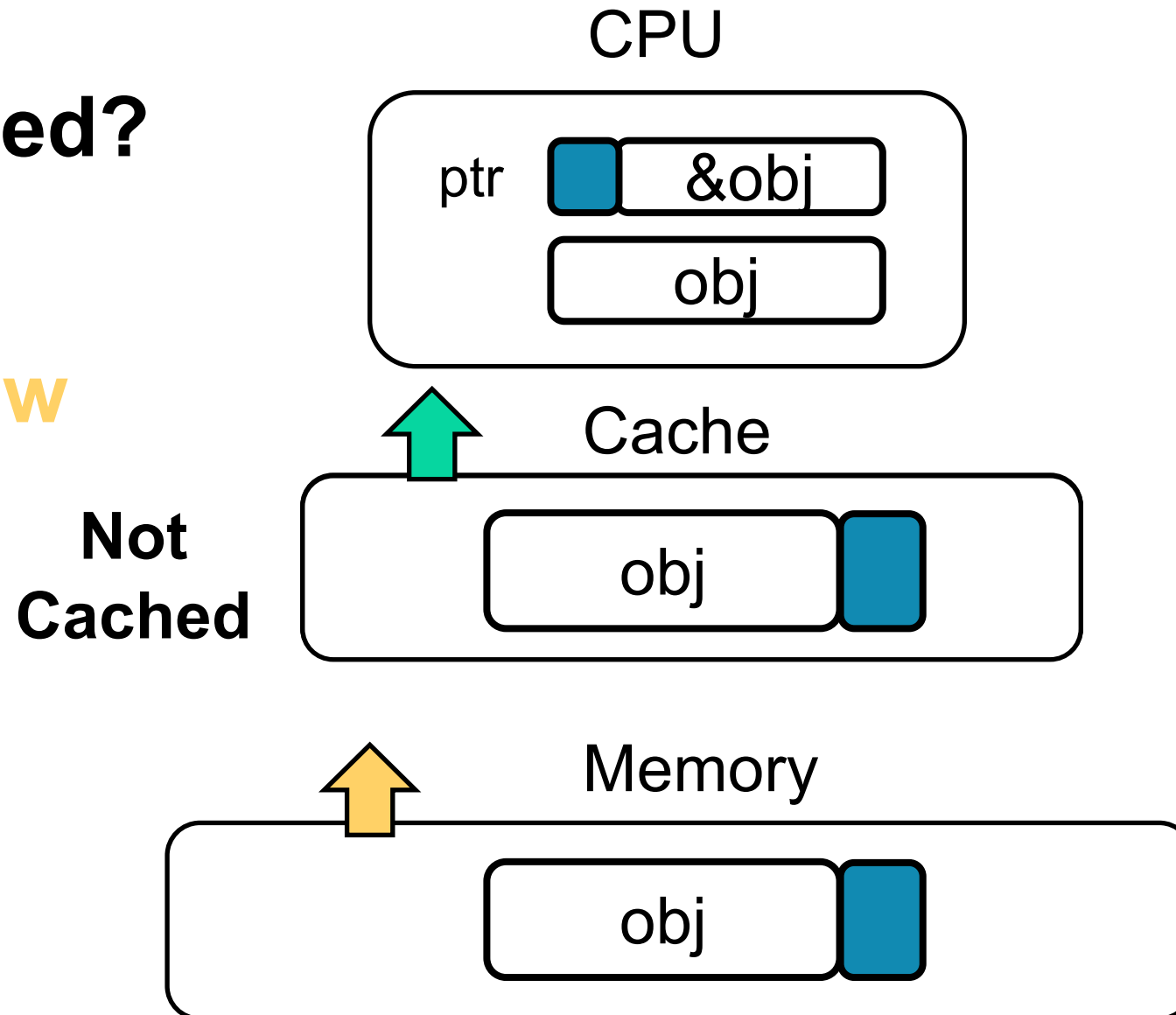
A. `ptr` has been accessed!

# What is Cache Side-Channel?

Q. Has obj been accessed?



A. ptr has NOT been accessed!



**Exploit cache side-channel**

**→ Leak whether an address is accessed**

# Roadmap

ARM Memory Tagging Extension

arm



Cache Side-Channel

Cache



Speculative Execution

if (cond)

False



Real-world MTE Bypass Attack



JS

MTE

MTE Tag Leakage Side-Channel

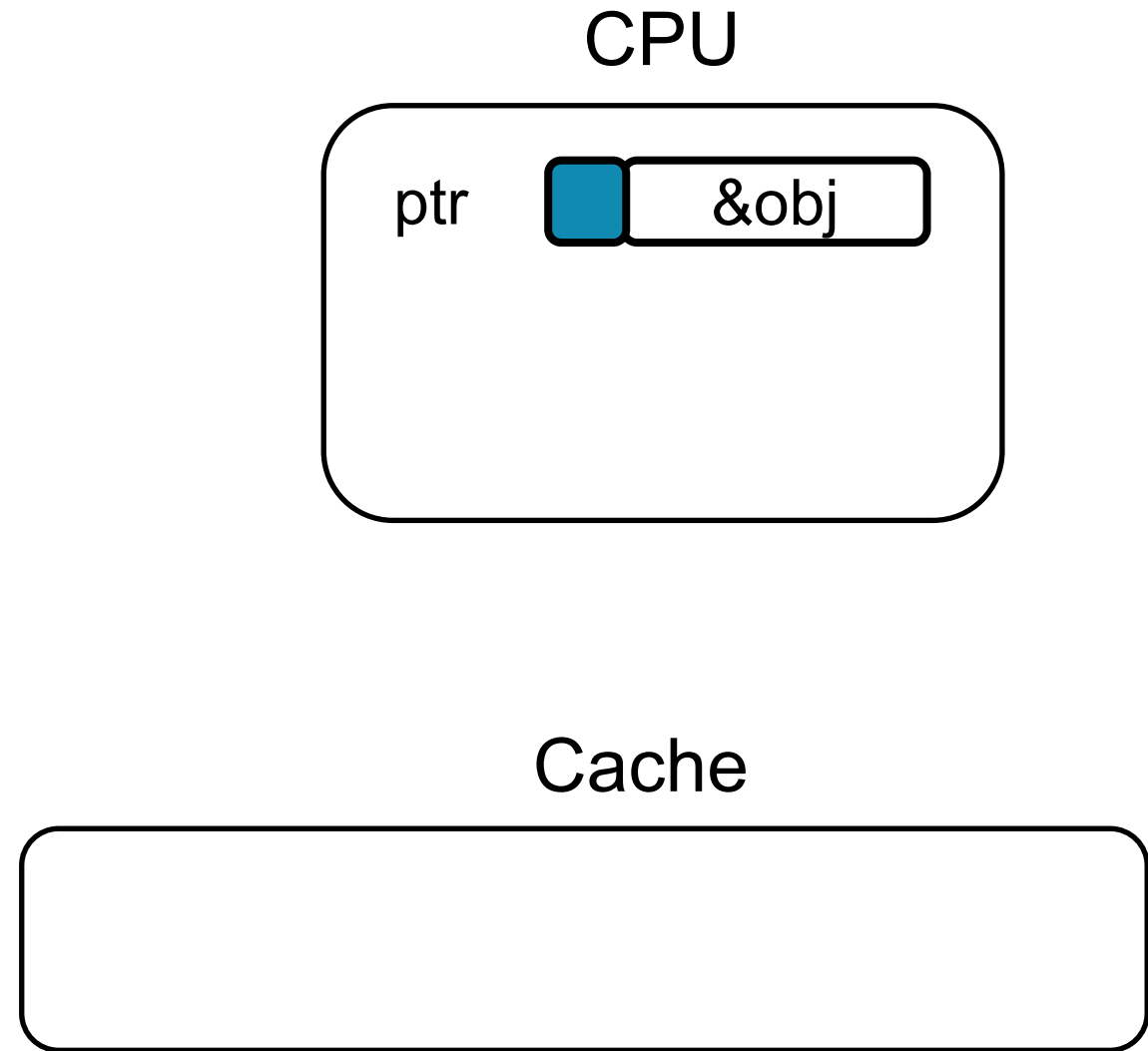


MTE

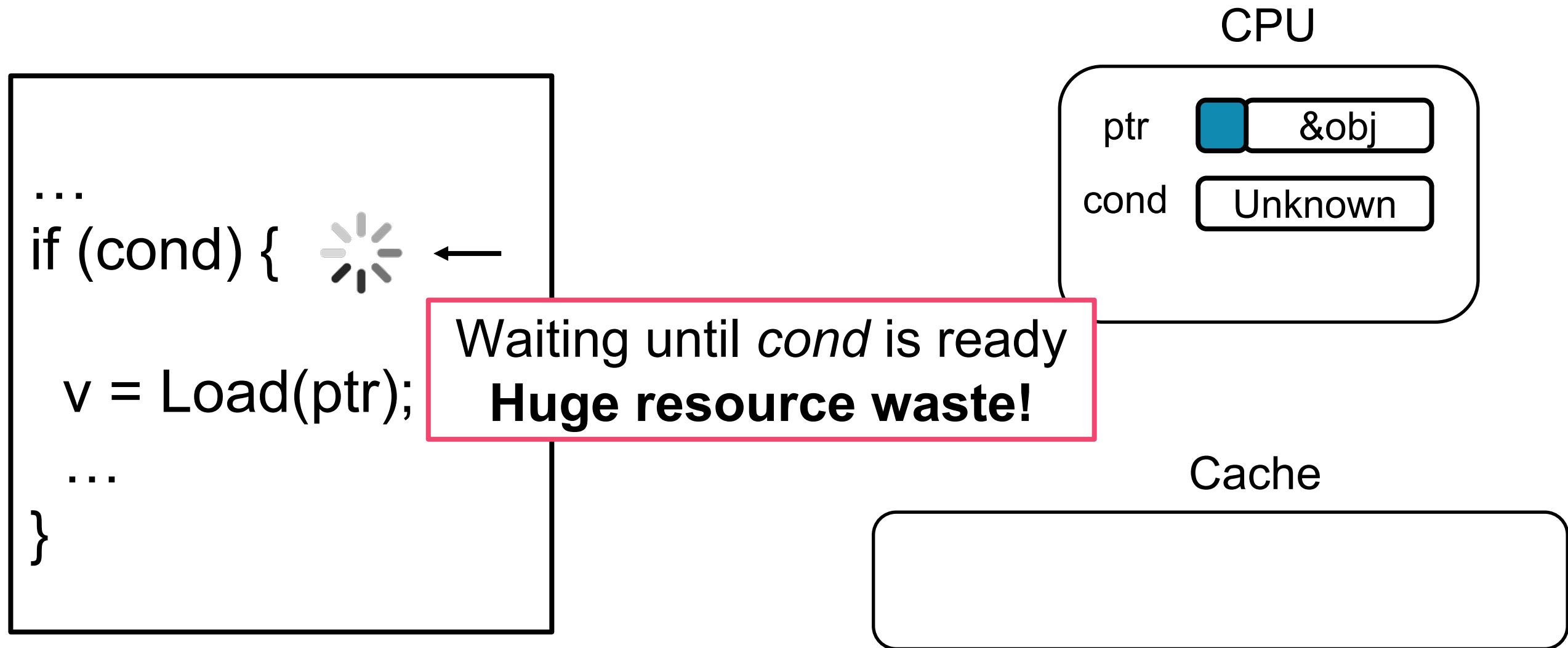


# What is Speculative Execution?


```
...  
if (cond) { ←  
  
    v = Load(ptr);  
  
    ...  
}
```

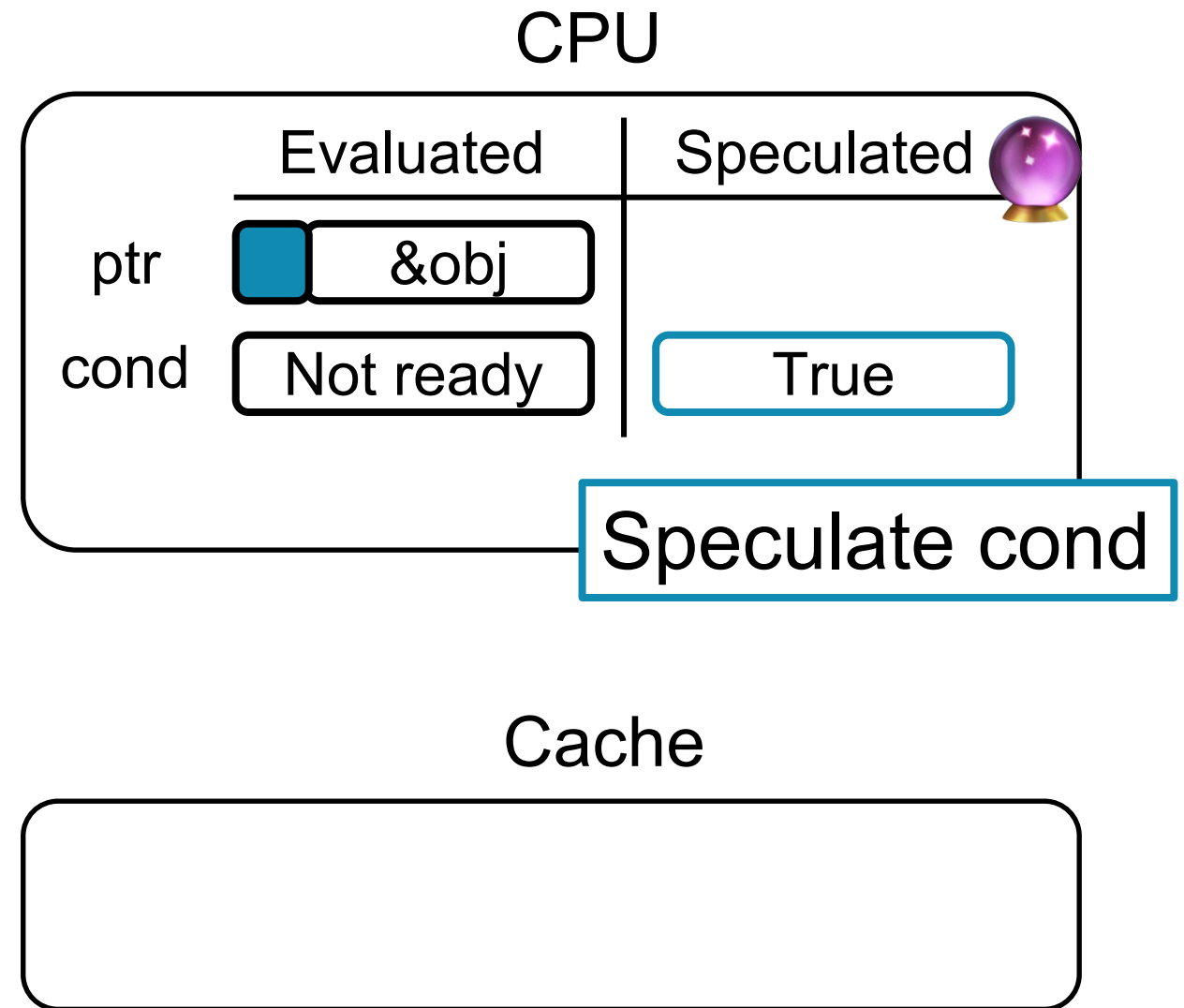


# What is Speculative Execution?





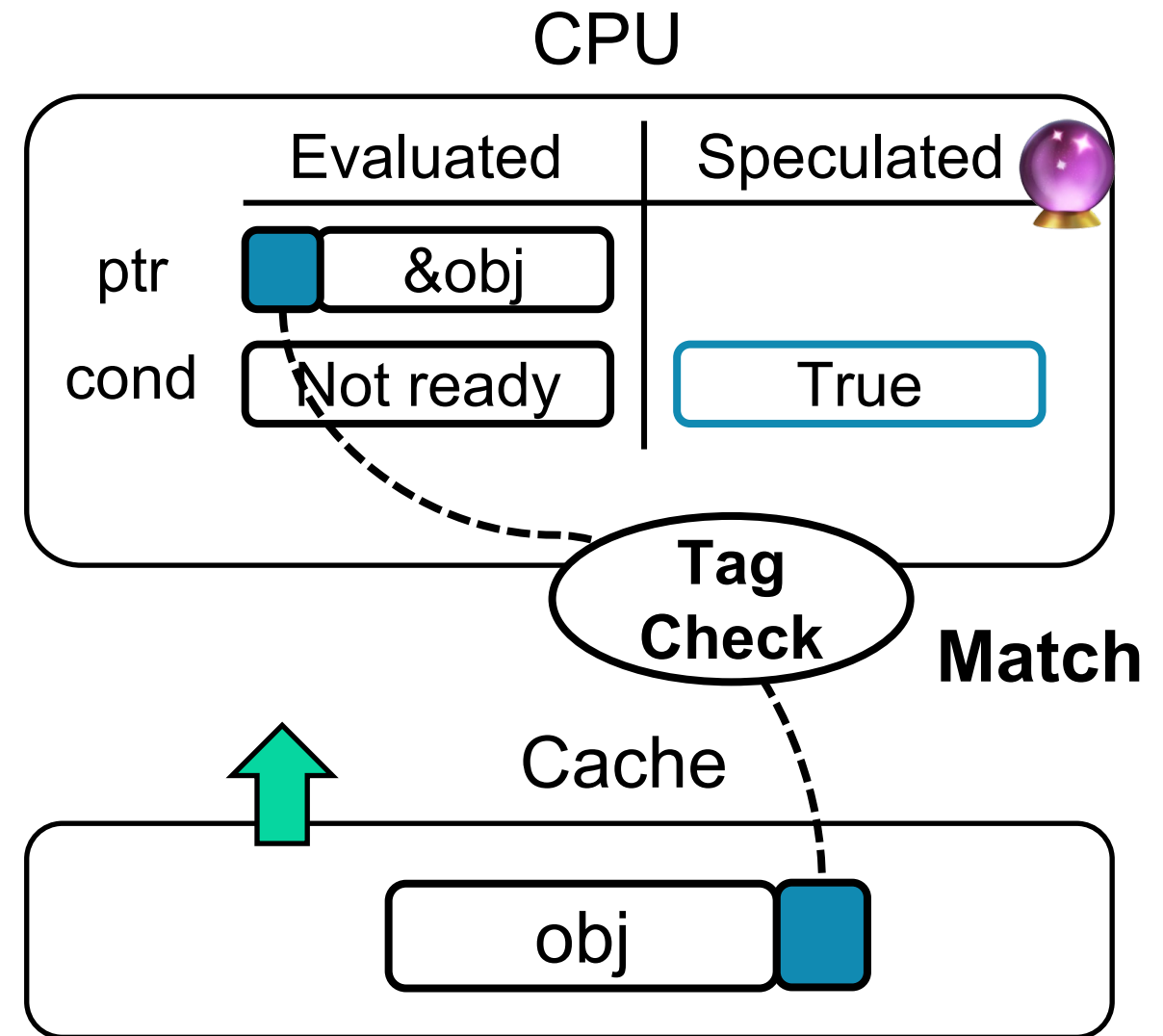
# What is Speculative Execution?

```
...  
if (cond) {  ←  
    v = Load(ptr);  
    ...  
}
```





# What is Speculative Execution?

```
...  
if (cond) {   
    v = Load(ptr);   
    ...  
}
```

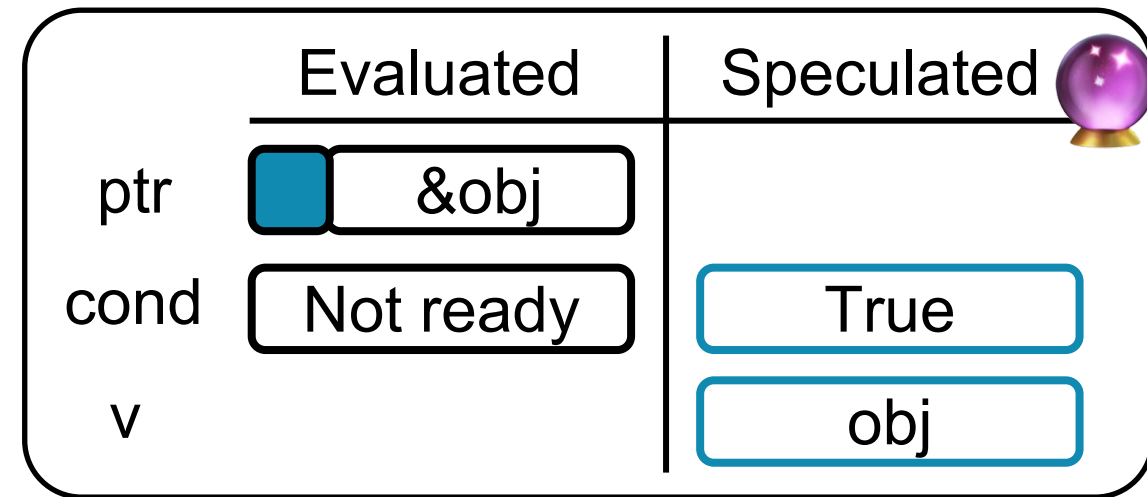




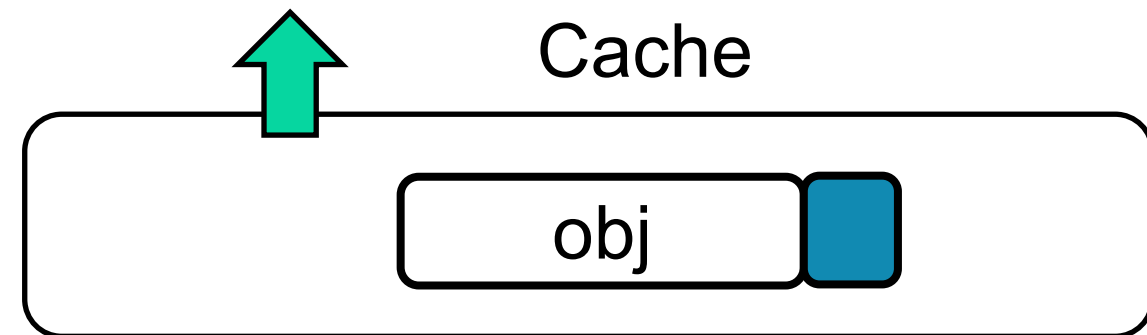
# What is Speculative Execution?

```
...  
if (cond) {   
    v = Load(ptr);   
    ...  
}
```

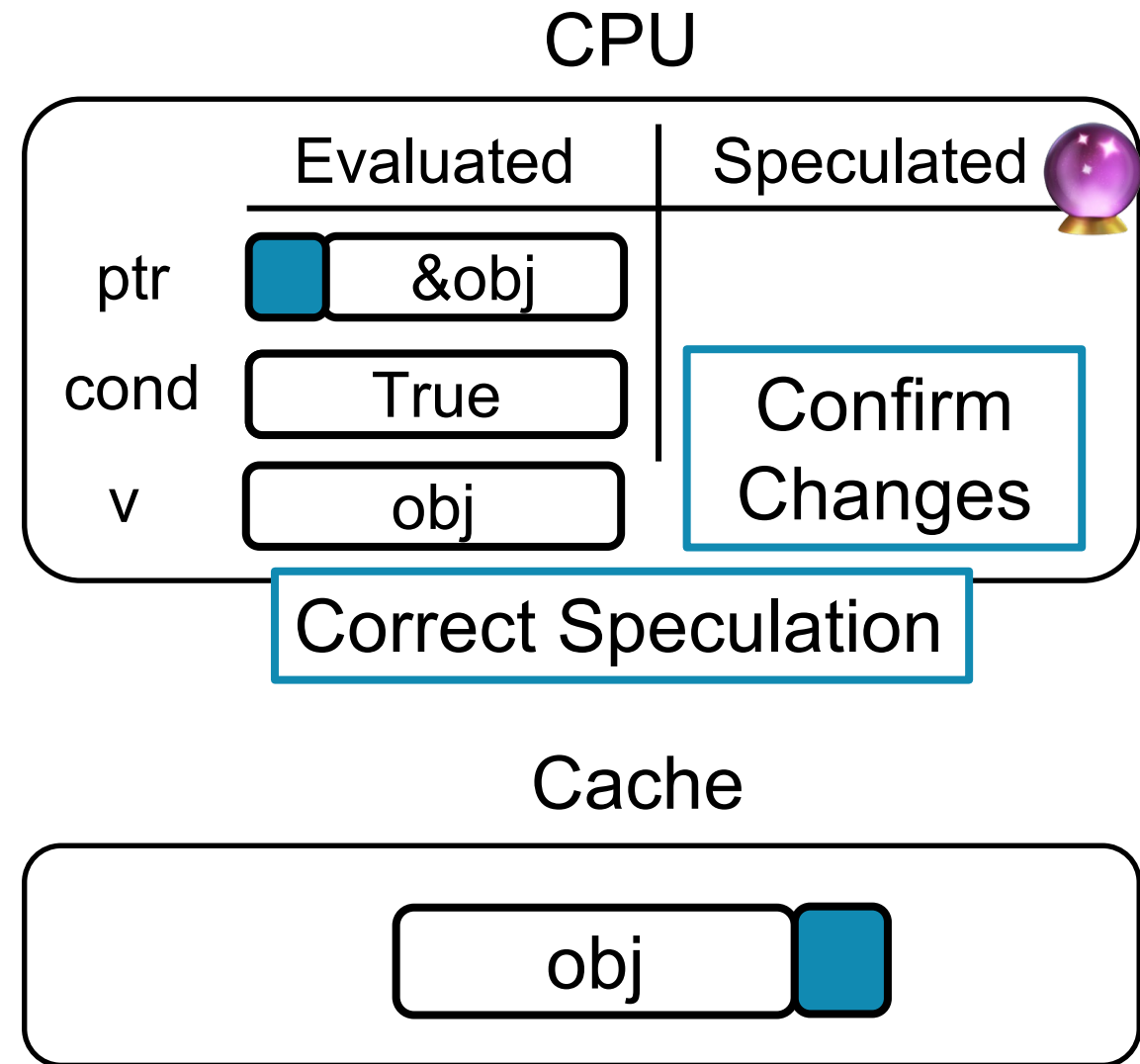
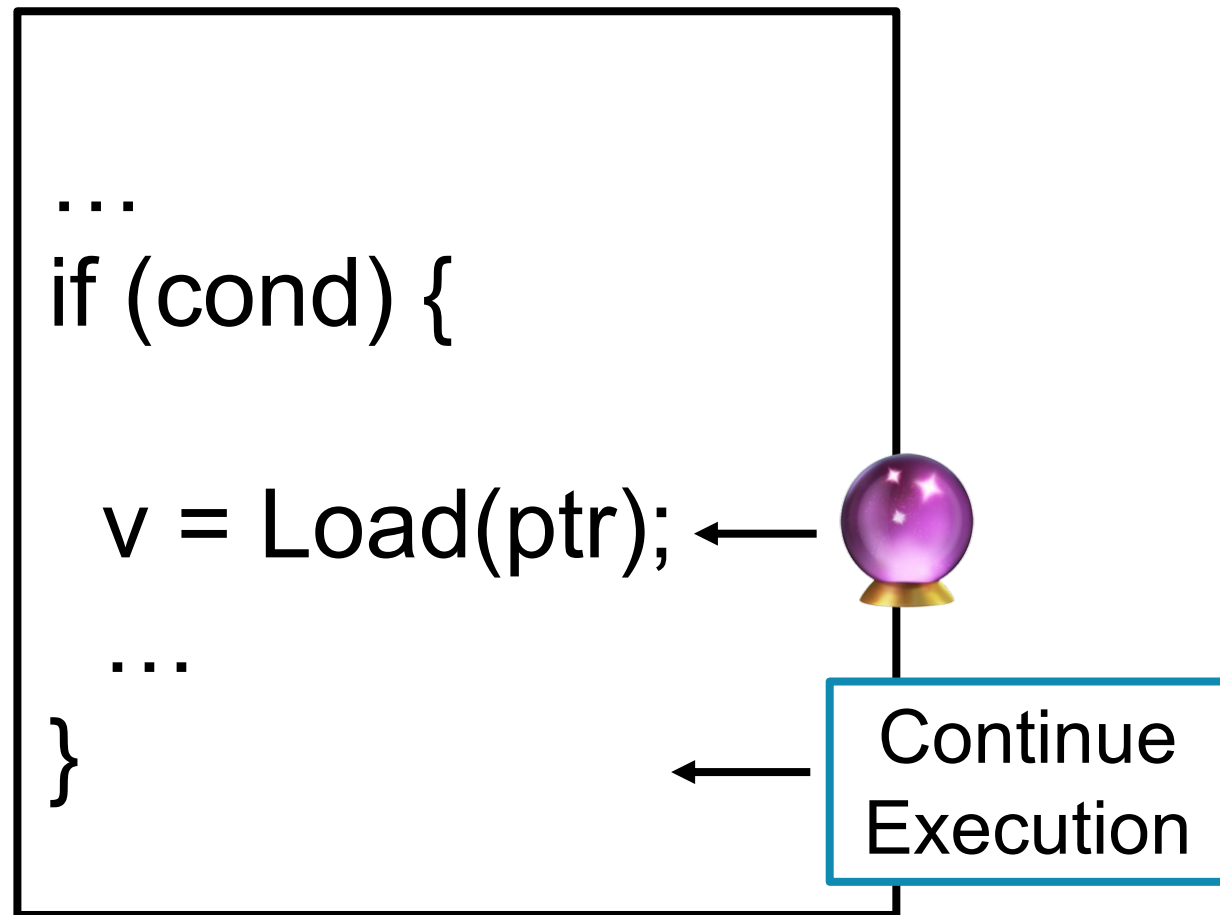
CPU



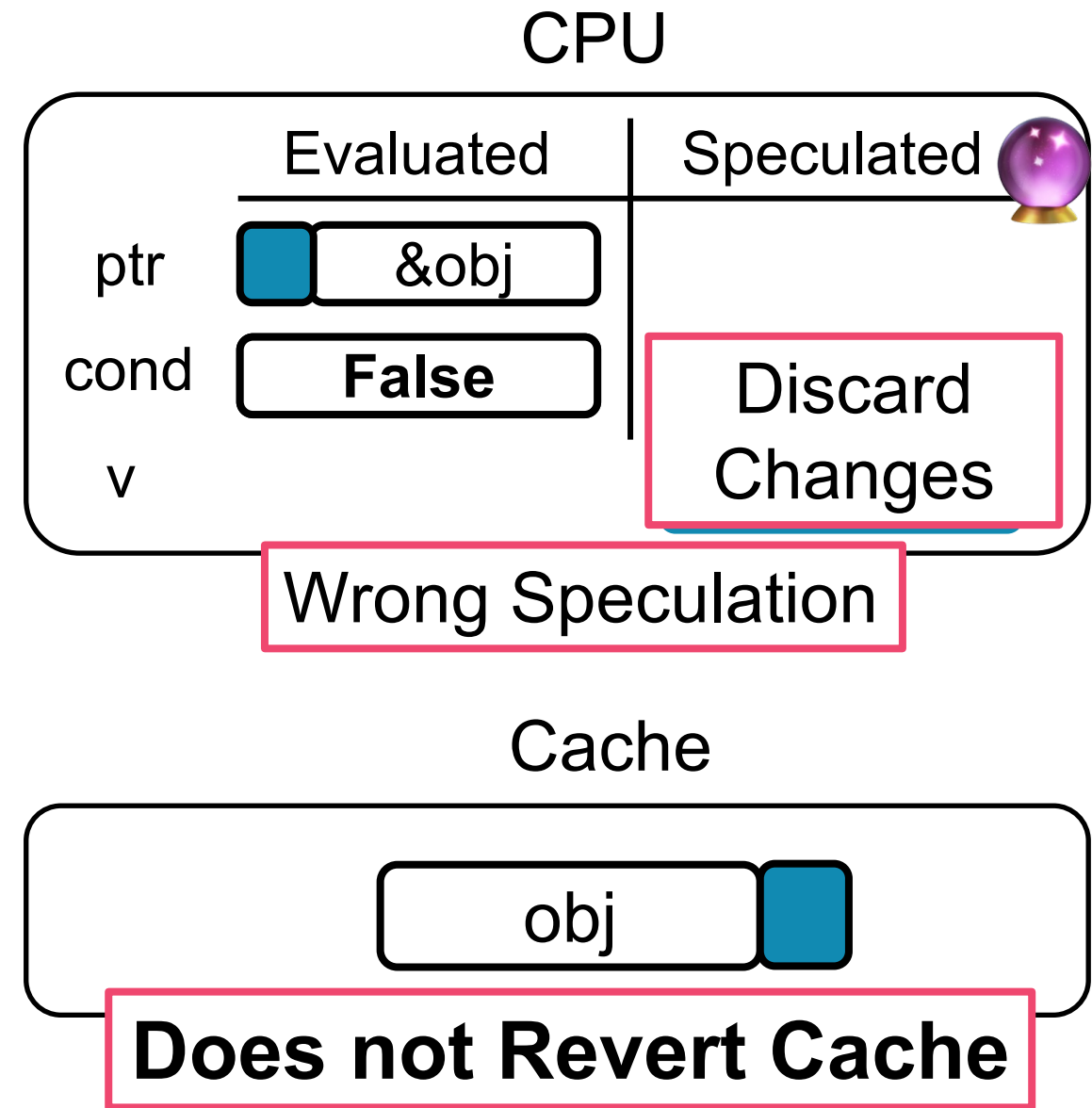
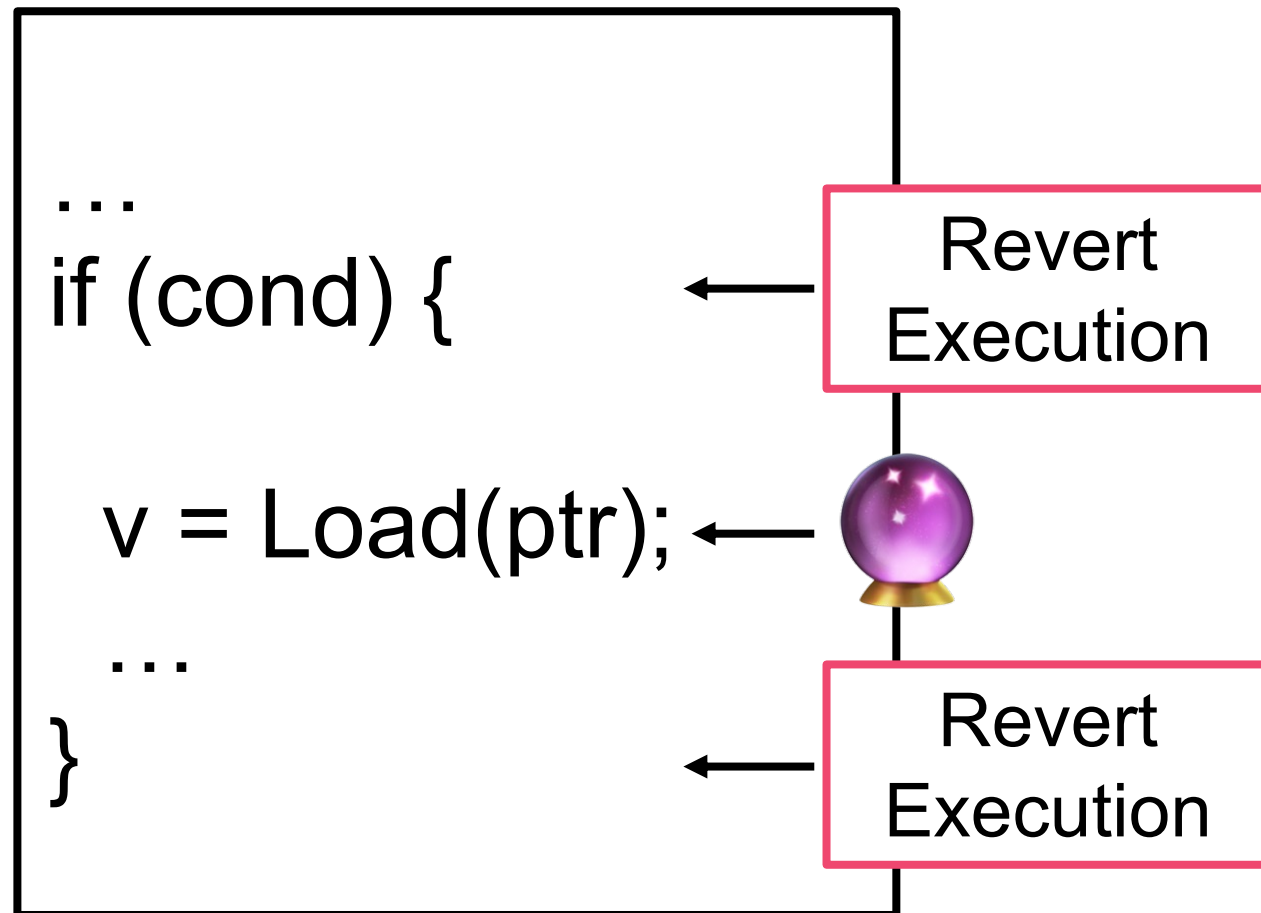
Cache



# What is Speculative Execution?

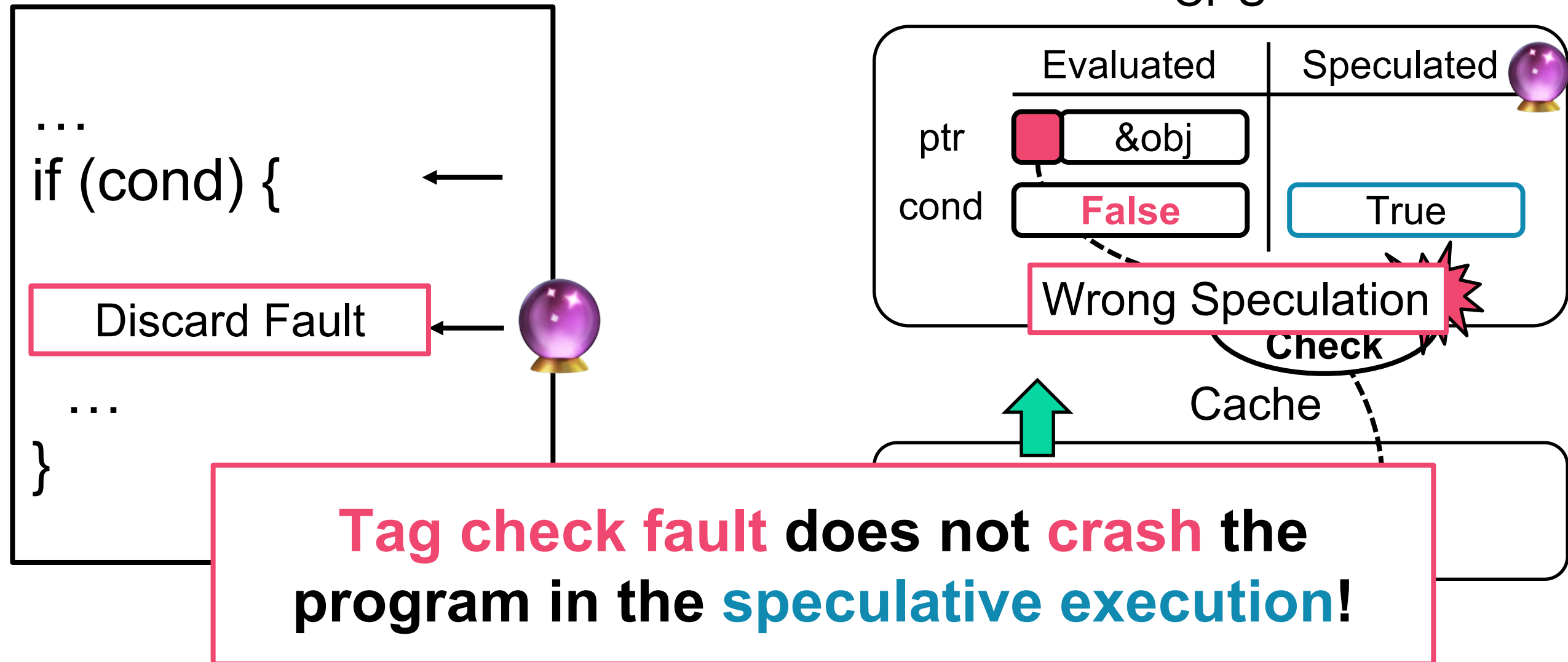


# What is Speculative Execution?



# Tag check fault on Speculative Execution?

CPU



**Exploit cache side-channel**

→ **Leak whether an address is accessed**

**Exploit speculative execution**

→ **Avoid crash on tag check fault**

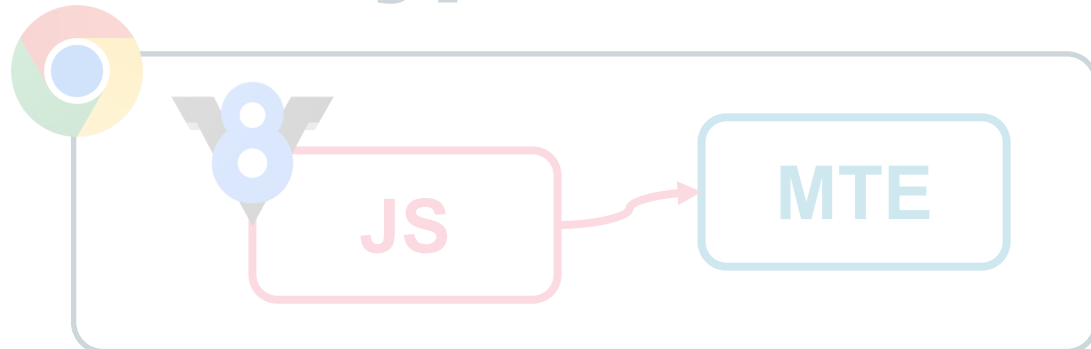
# Roadmap

ARM Memory Tagging Extension

arm



Real-world MTE Bypass Attack



Cache Side-Channel

Cache



Speculative Execution

if (cond)

True False



MTE Tag Leakage Side-Channel

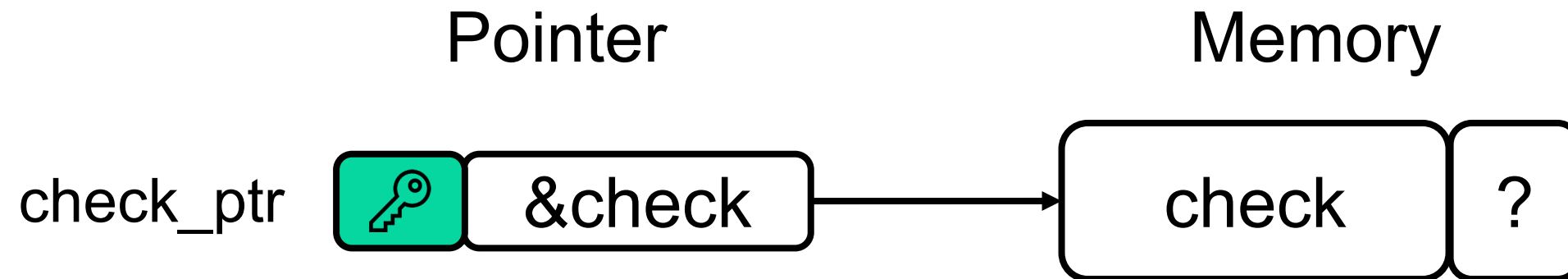


MTE



# MTE Side-channel attack

**Goal: Leak the memory tag given a pointer**



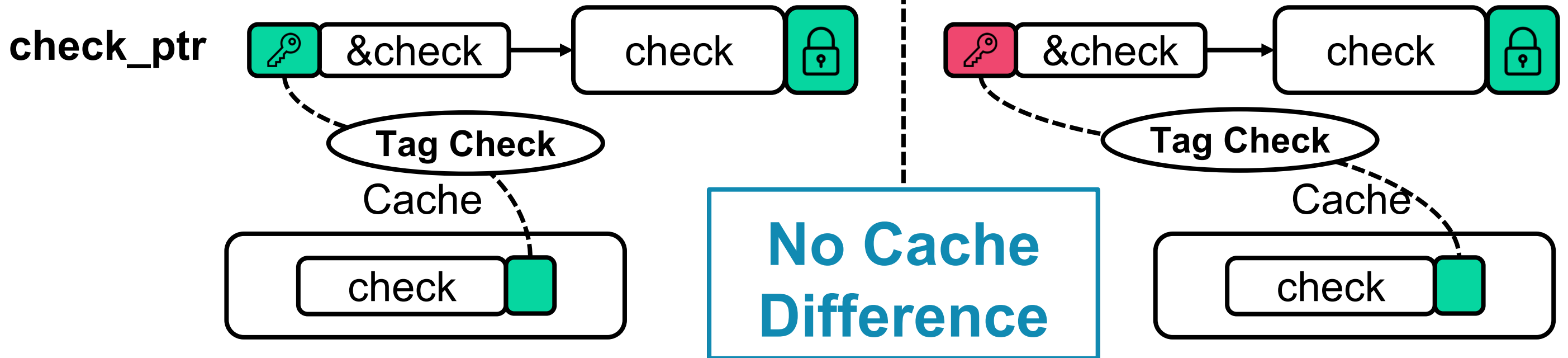
# MTE Side-channel attack

Two test cases:

```
Access(check_ptr);
```

A. **Valid** tag in check\_ptr

B. **Invalid** tag in check\_ptr





# MTE Side-channel attack

Two test cases:

```
Access(check_ptr); Access(test_ptr);
```

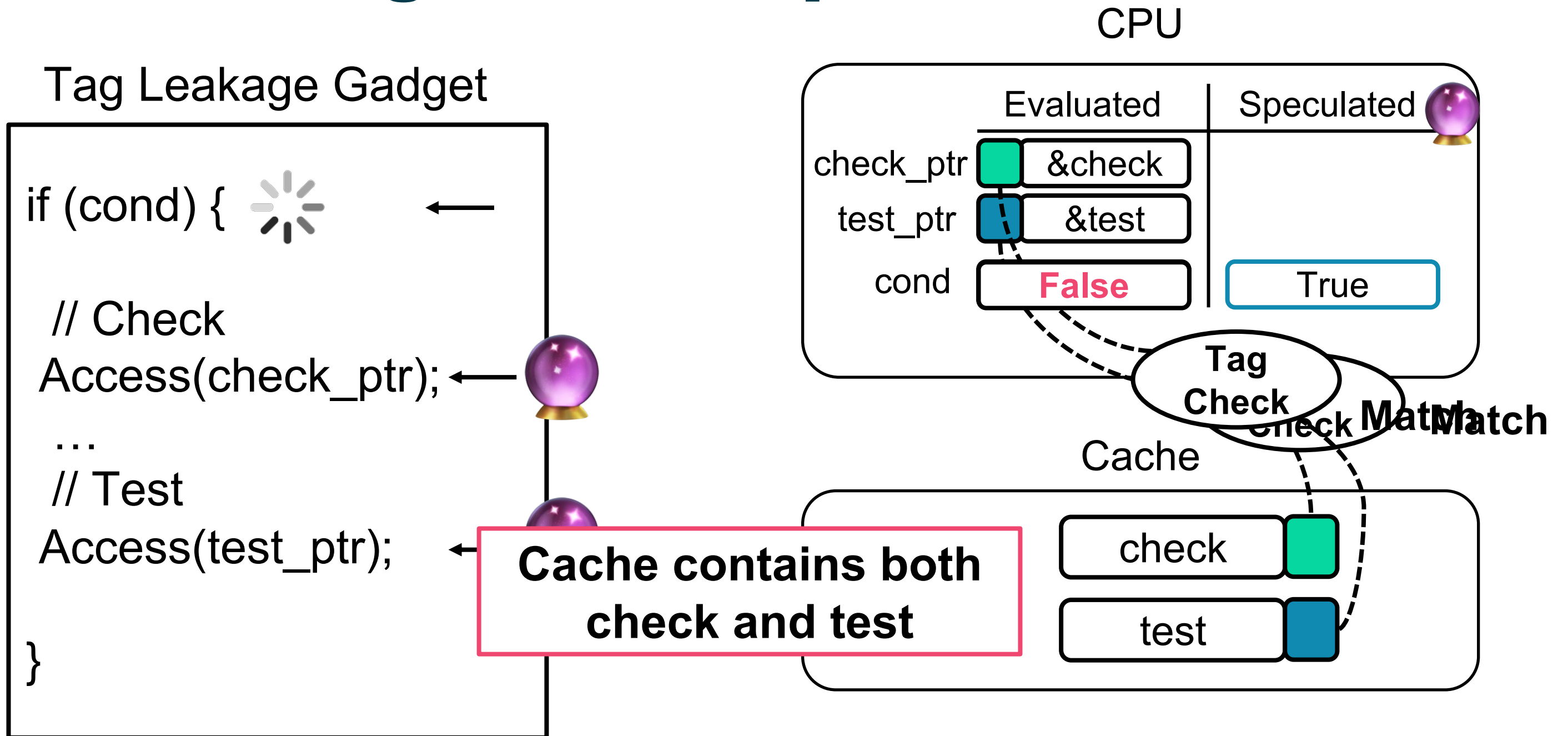
A. **Valid** tag in check\_ptr

B. **Invalid** tag in check\_ptr



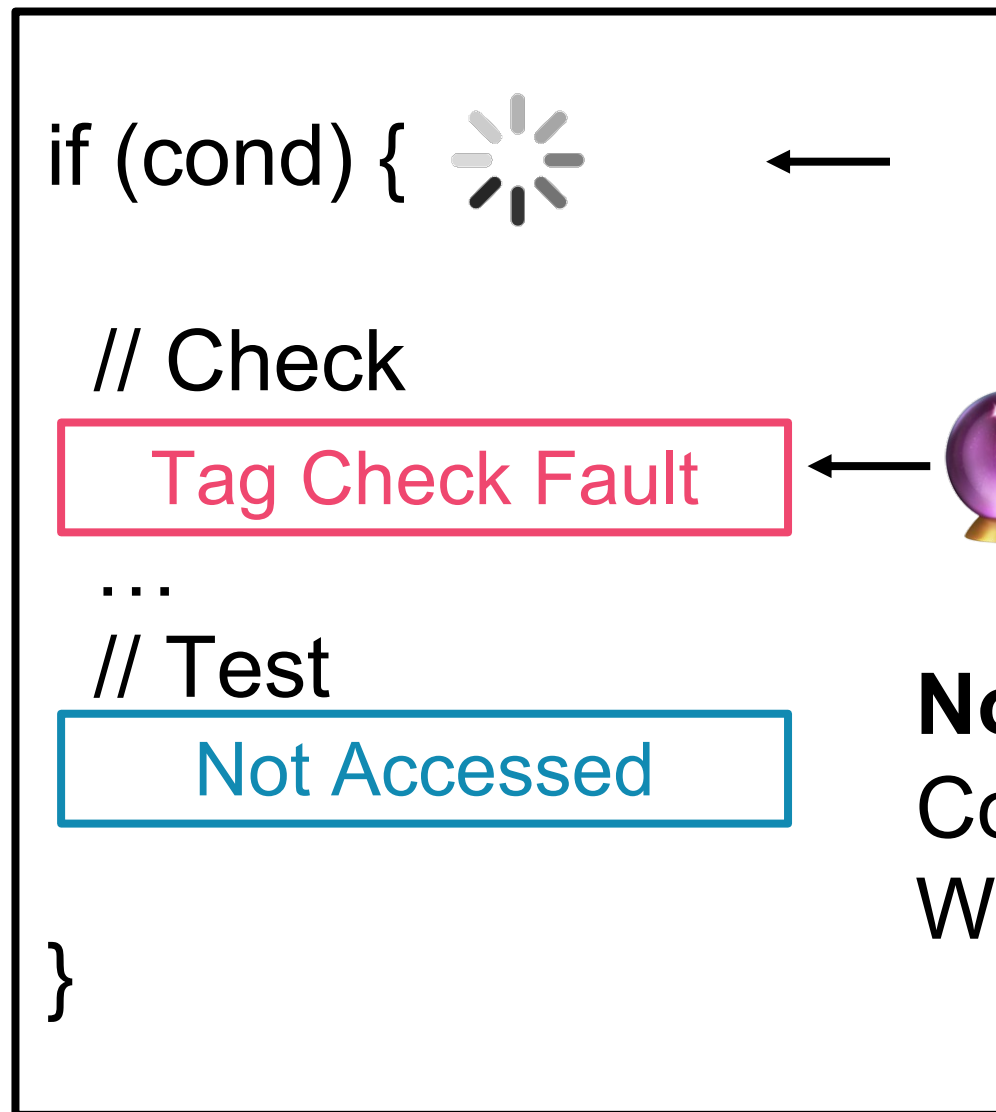
Cache  
Difference?

# A. Valid tag in check\_ptr

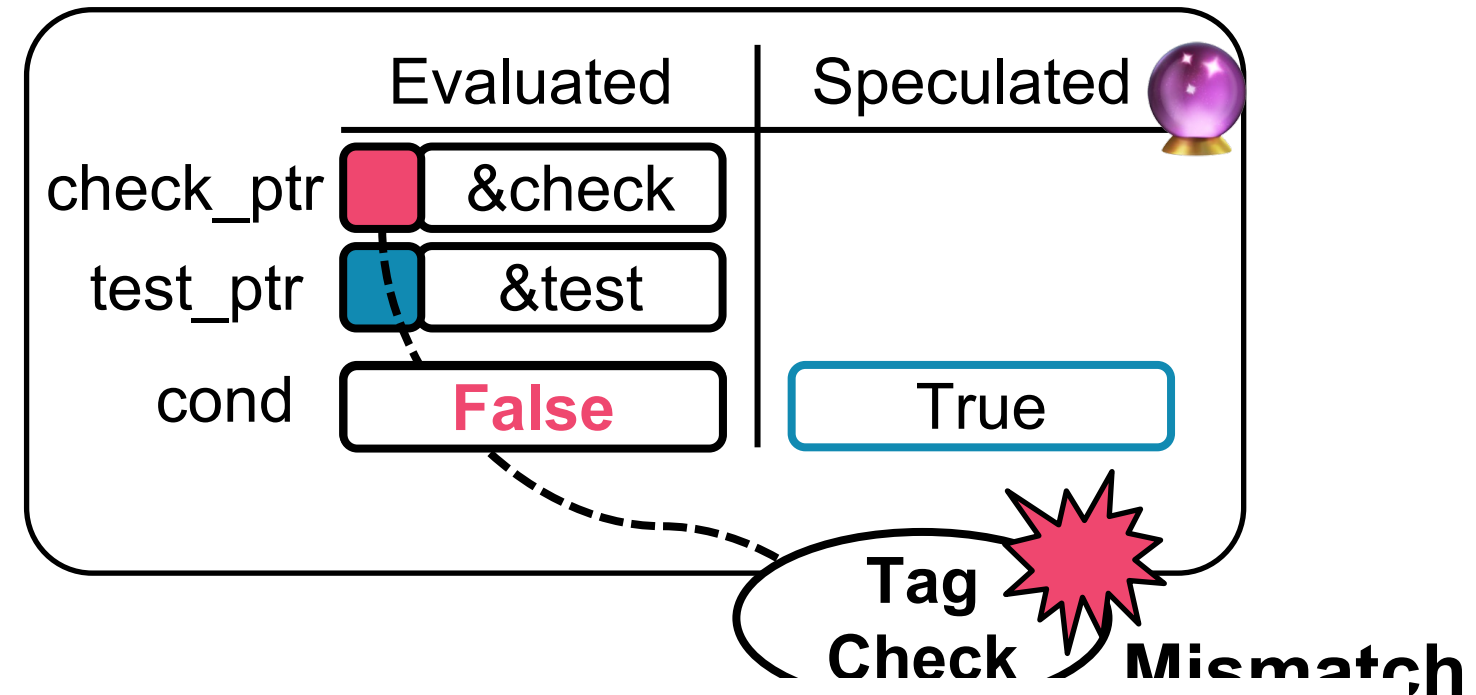


# B. Invalid tag in check\_ptr

## Tag Leakage Gadget



CPU



**No reason to continue speculative execution**

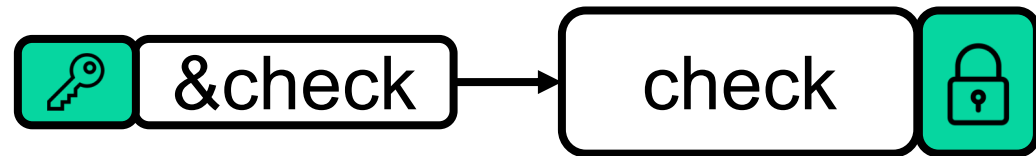
Correct spec → (synchronous) tag check fault

Wrong spec → Revert execution

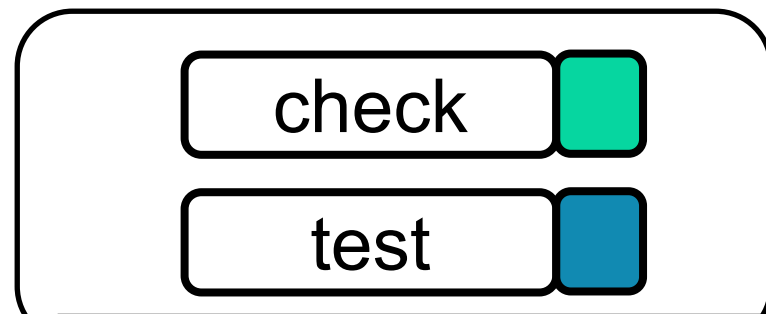
**check, not test**

# Leak by Cache Side-Channel

## A. **Valid** tag in check\_ptr



Cache

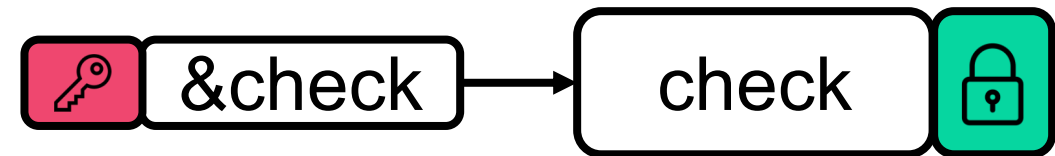


```
Load(test_ptr);
```

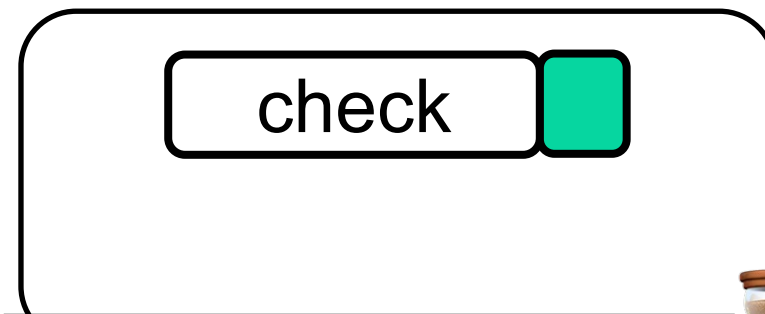


**Fast**

## B. **Invalid** tag in check\_ptr



Cache



```
Load(test_ptr);
```



**Slow**

**Leak whether the tag is **Valid/Invalid** by **test\_ptr** access latency!**

# Do new MTE chips contain the tag leakage side-channels?

## PACMAN – ISCA 2022, DEF CON 30

- Discovered a [Pointer Authentication Code \(PAC\)](#) side-channel

## MTE as Tested – Google Project Zero, POC 2023

- Attempted to find a [MTE tag](#) side-channel → **Failed**

```
if (cond) {  
    // Check
```

## Our work

- **Found 2 Tag Leakage Gadgets + Suspected Root Causes**
- Gadget poc: <https://github.com/compsec-snu/tiktag>
- Detailed analysis in our paper: <https://arxiv.org/abs/2406.08719>

```
    :k_ptr;
```

## StickyTags – VUSec, IEEE S&P 2024

- Orthogonally found one of our tag leakage gadgets

```
    Access(test_ptr);  
}
```

# Gadget 1: Multiple Loads

```
if (cond) {  
    ...  
    // Check: 2+ load  
    *check_ptr; Tag Check Fault  
    *check_ptr; Tag Check Fault  
    ...  
    // Test: load/store  
    *test_ptr; No Access  
}
```

Suspected root cause

- On **multiple faults**, the CPU **re-speculates that the speculation was wrong**  
=> **stop/reduce speculations** in branch speculation and memory prefetcher

(12) **United States Patent**  
Cai et al.

(10) **Patent No.:** US 11,526,356 B2  
(45) **Date of Patent:** Dec. 13, 2022

(54) **PREFETCH MECHANISM FOR A CACHE**  
STRUCTURE

(56)

**References Cited**

U.S. PATENT DOCUMENTS

*The wrong path event... provides a hint that the processor pipeline may have fetched one or more instructions that do not require execution. ... some examples are invalid memory accesses, ...*

# Gadget 2: Store-to-Load Forwarding

```

if (cond) {
  // Check: store-to-load
  *check_ptr = val;
  val = *check_ptr;

  // Test: dependent load/store
  *(test_ptr+val);
}

```

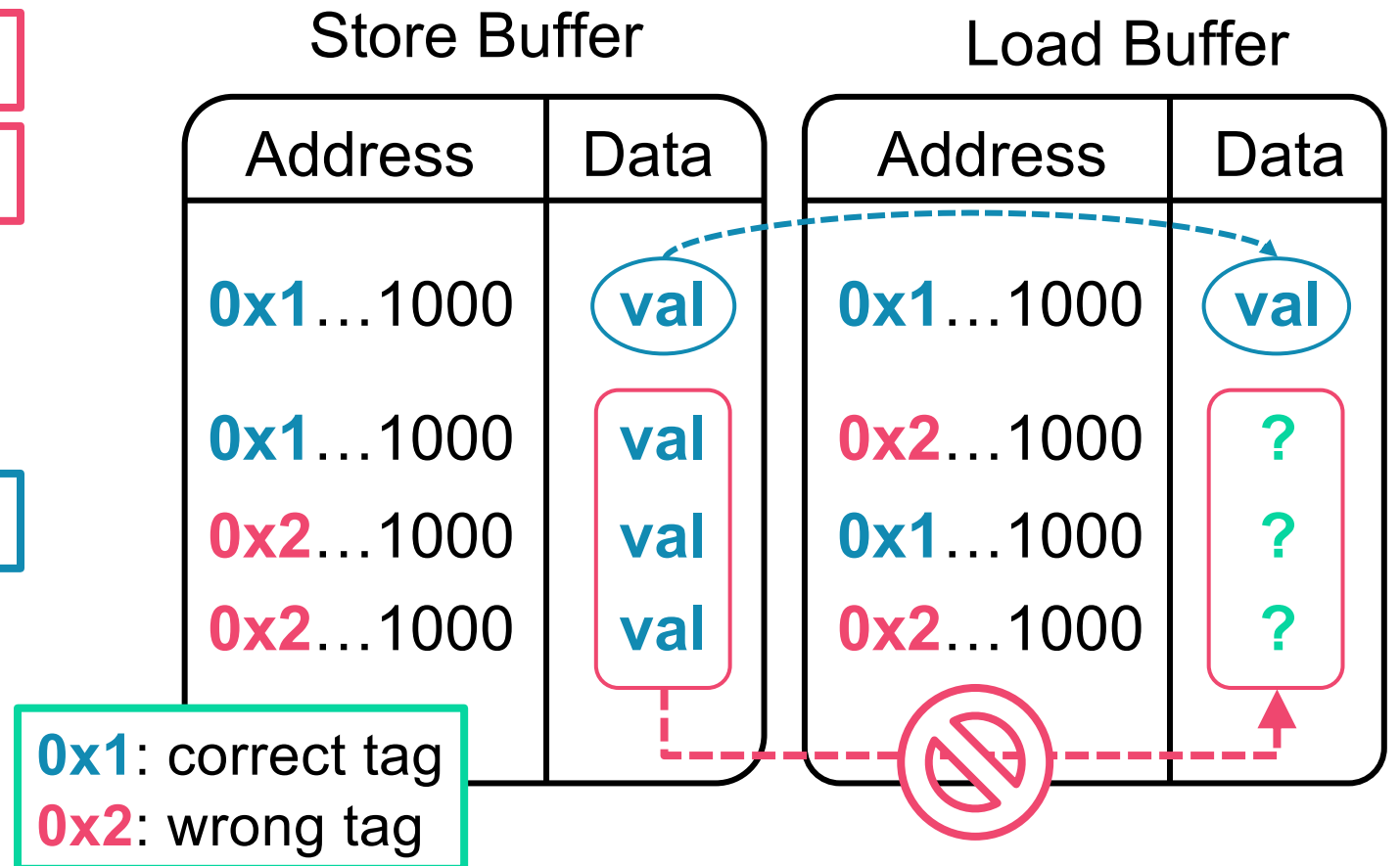
Tag Check Fault

Tag Check Fault

No Access

Suspected root cause

- On **tag check fault**, the CPU blocks *store-to-load forwarding*



# Roadmap

ARM Memory Tagging Extension

arm



Cache Side-Channel

Cache



Speculative Execution

if (cond)  
True ↙ False ↘



Real-world  
MTE Bypass Attack



JS

MTE

MTE Tag Leakage  
Side-Channel





# Real-world MTE-Enabled Software

- MTE became recently available
- Software systems that provide (optional) MTE support



**android**

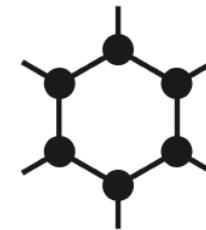


**Google  
Chrome**



**Linux  
Kernel**

## Secure OSes



**GrapheneOS**



**Unikraft**



**OPTEE**

- More software systems are likely to adopt MTE in the near future

# Real-world Gadgets & Attacks

## When MTE is enabled

### 1. Google Chrome V8 Engine

Constructed exploitable Gadget 2 from JavaScript

→ Leak MTE tag of the renderer memory

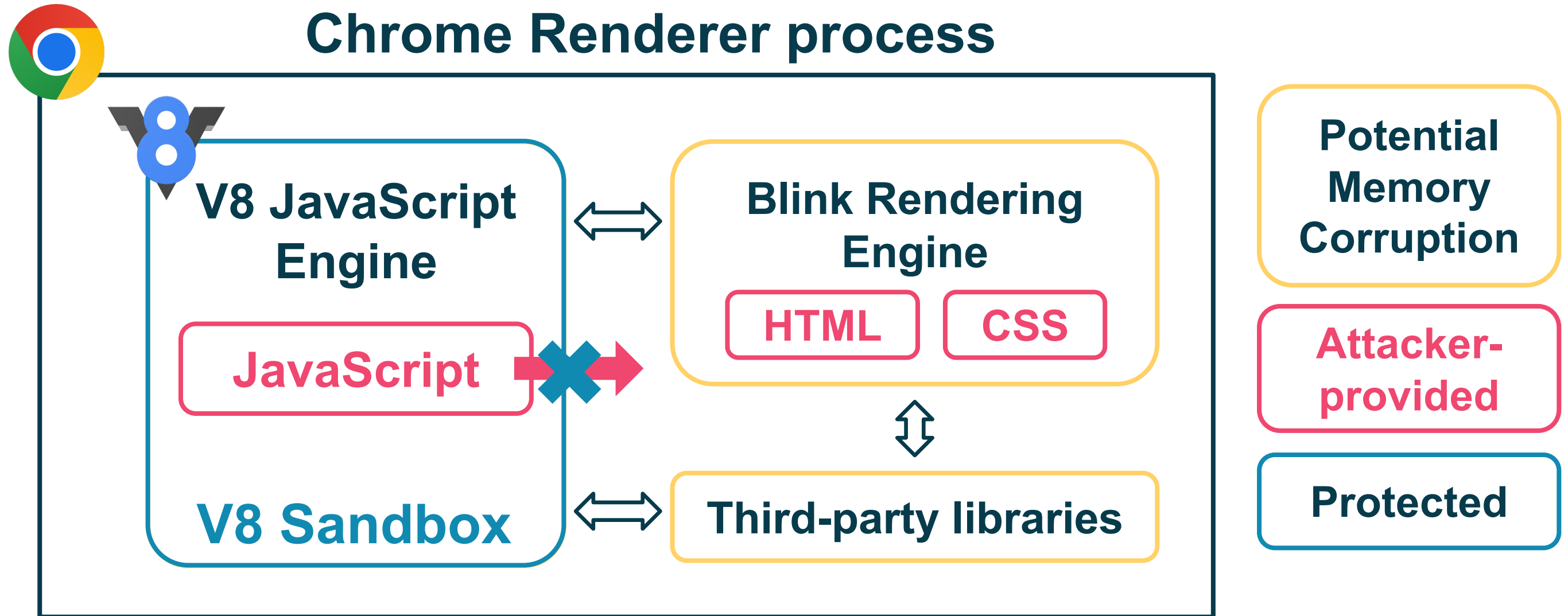
### 2. Linux kernel

Found potential Gadget 1 in `snd_timer()`

→ Leak MTE tag of the kernel memory from user space

Refer to our paper for the details: <https://arxiv.org/abs/2406.08719>

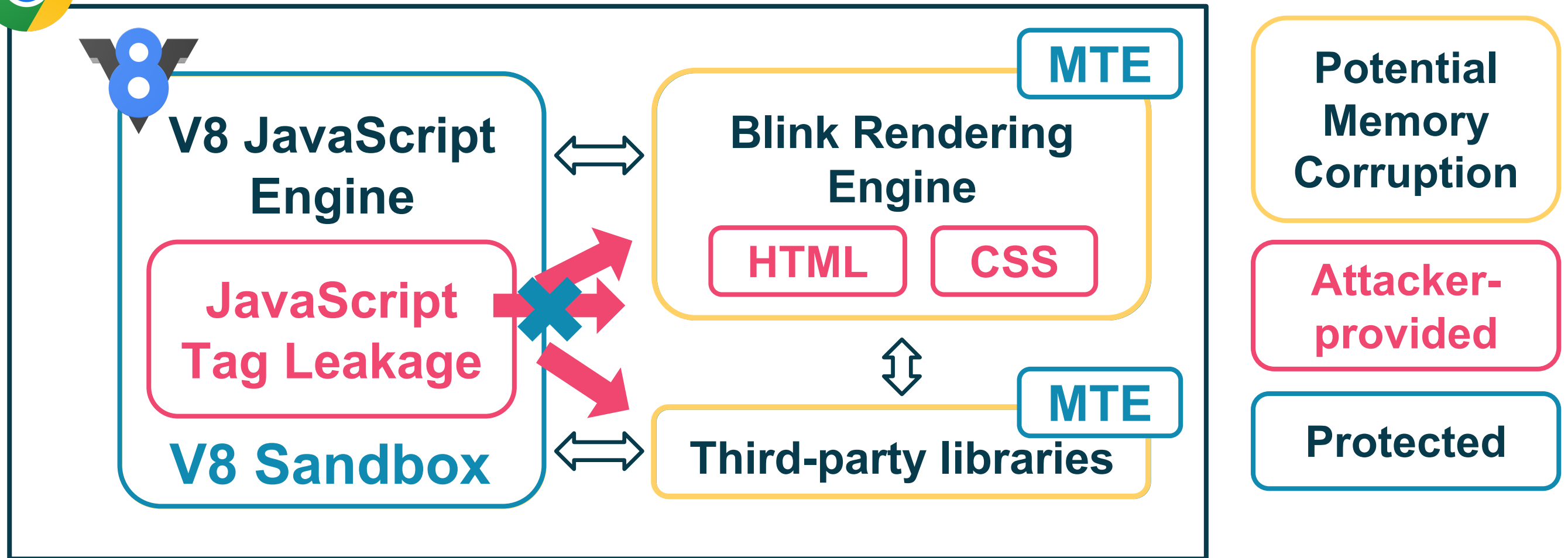
# Google Chrome Threat Model



# Google Chrome Threat Model




## Chrome Renderer process



# Gadget 2 from JavaScript

```
if (cond) {  
    check[idx] = val;  
    val = check[idx];  
    x = test[val];  
}
```

**Speculative Execution**



**idx**: out-of-bounds index (64-bit)

**check[idx]**: check\_ptr

**test[val]**: test\_ptr

# Gadget 2 in Google V8 (JavaScript)

```
TagLeak(target) {  
  for (let tag=0; tag < 16; ++tag) {  
    idx = AddrToldx(tag, target);  
    if (cond) {  
      check[idx] = val;  
      val = check[idx];  
      x = test[val];  
    }  
    time[tag] = Measure(test[val]);  
  }  
  return time.indexOf(min(time));  
}
```

for (let tag=0; tag < 16; ++tag) { ← Iterate all tag values  
 idx = AddrToldx(tag, target); ← out-of-bounds index

```
if (cond) {  
  check[idx] = val;  
  val = check[idx];  
  x = test[val];  
}
```

**Tag Leakage Gadget**



Valid tag	Invalid tag
No fault	Tag Check Fault
No fault	Tag Check Fault
Access	No Access
<b>Fast</b>	<b>Slow</b>

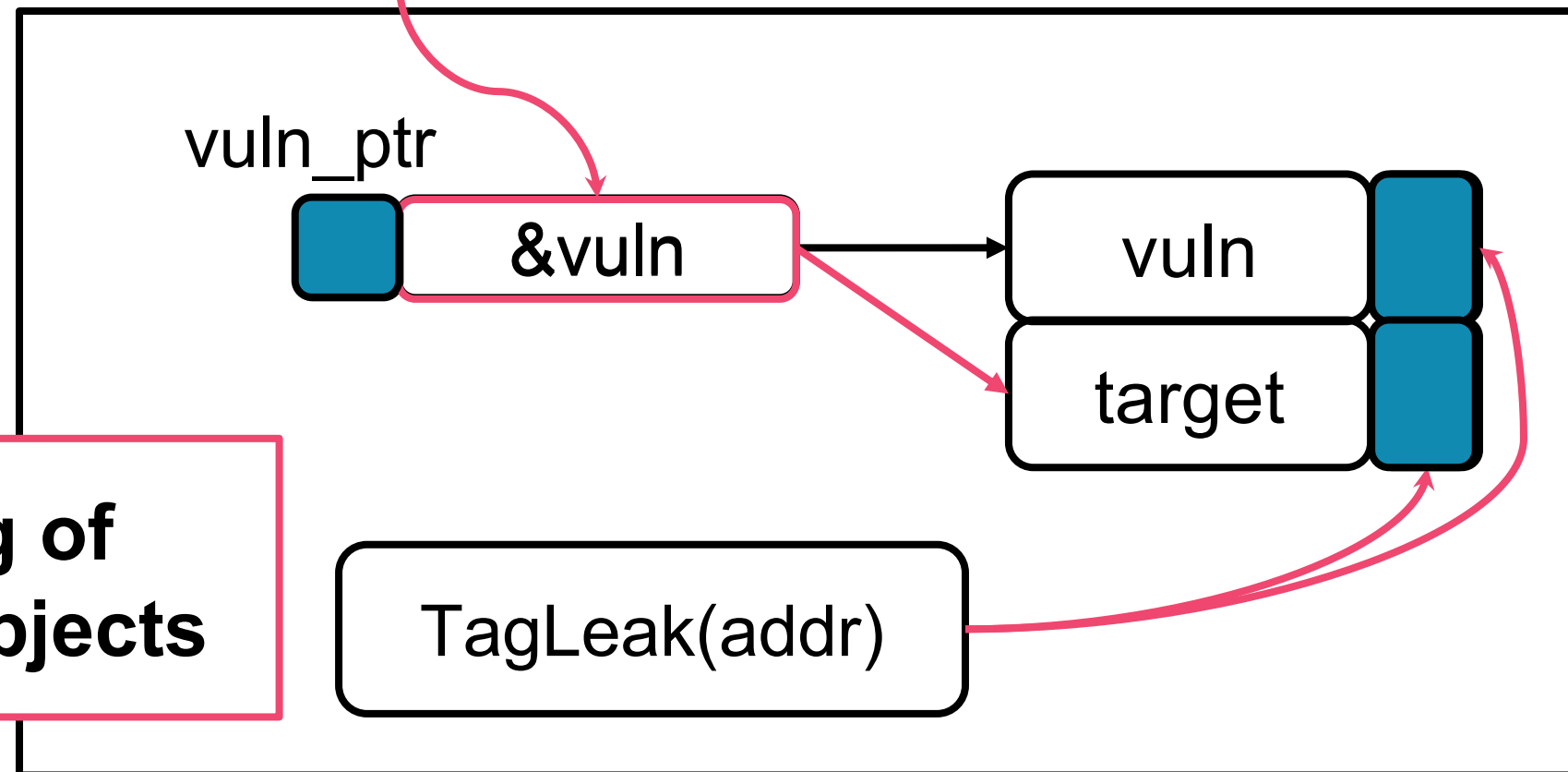
**Tag Leaked!**

# Chrome MTE Bypass Attack

Trigger memory corruption if tag match is expected




Chrome

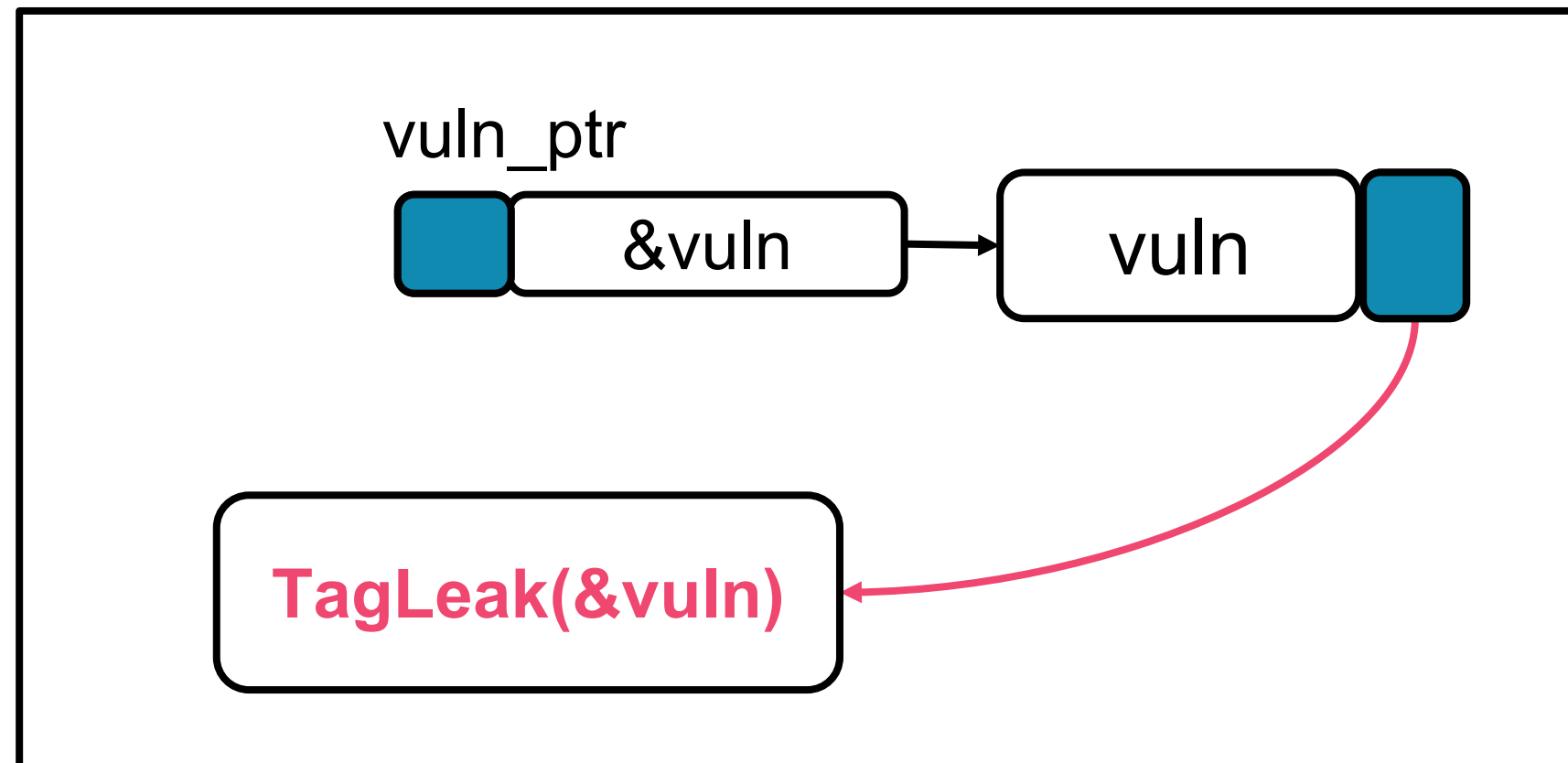


Leak tag of memory objects

# 1. Leak MTE Tag of vulnerable object



vuln.tag = 



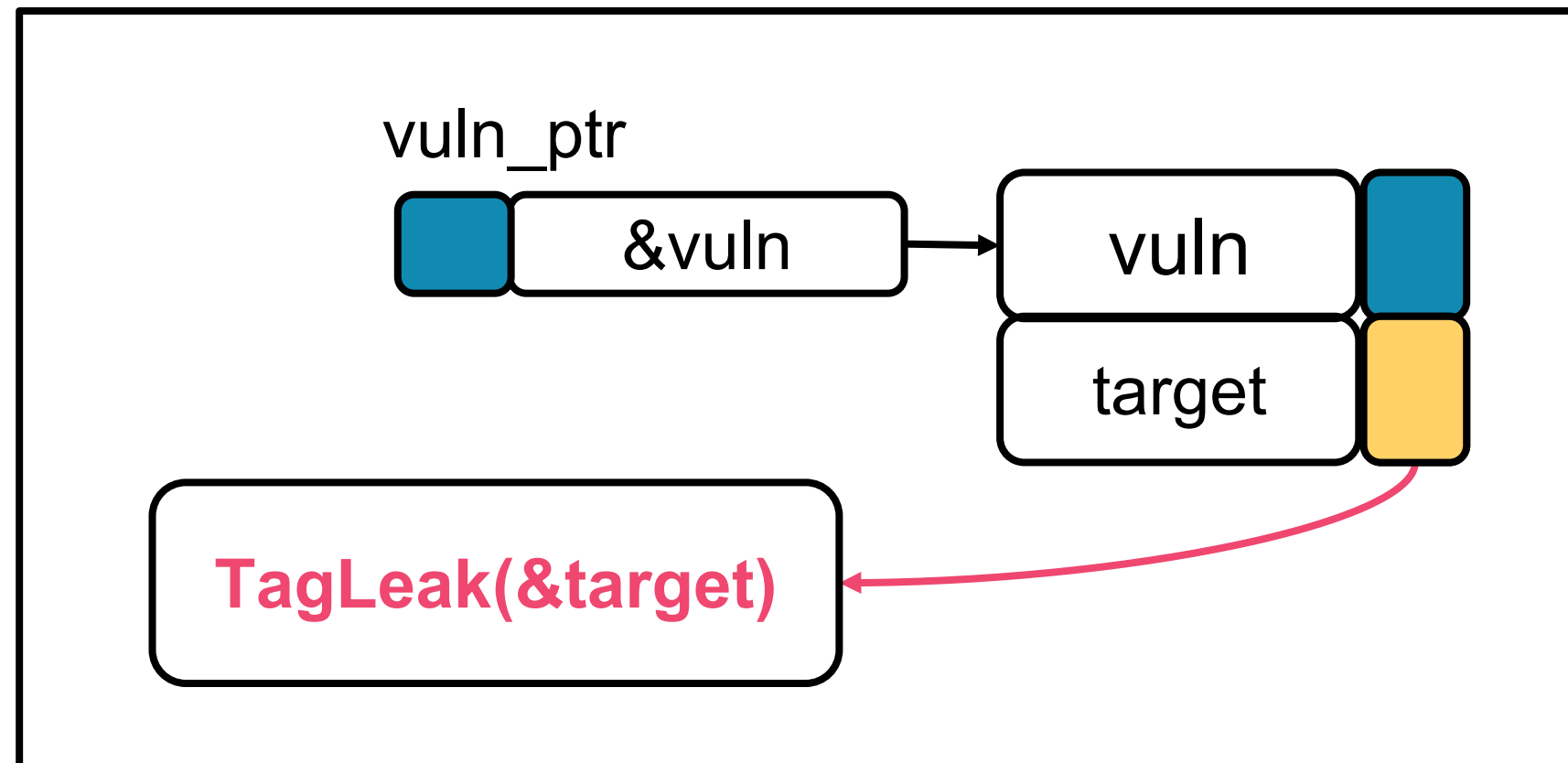


## 2. Leak MTE Tag of target object



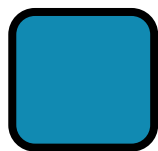
vuln.tag = 

target.tag = 



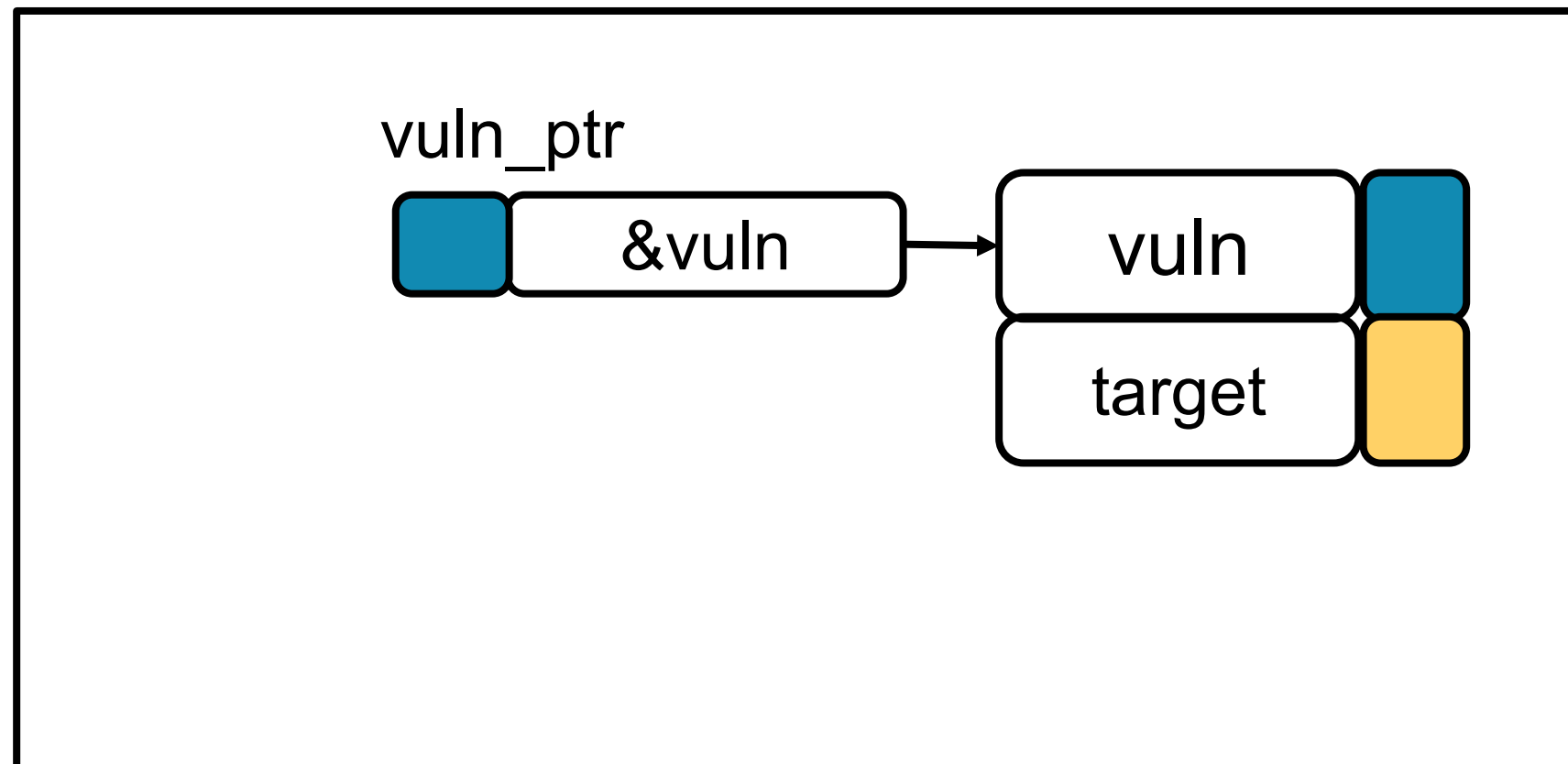
# 3. Reallocate target on tag mismatch



vuln.tag = 


target.tag = 

**vuln.tag != target.tag**

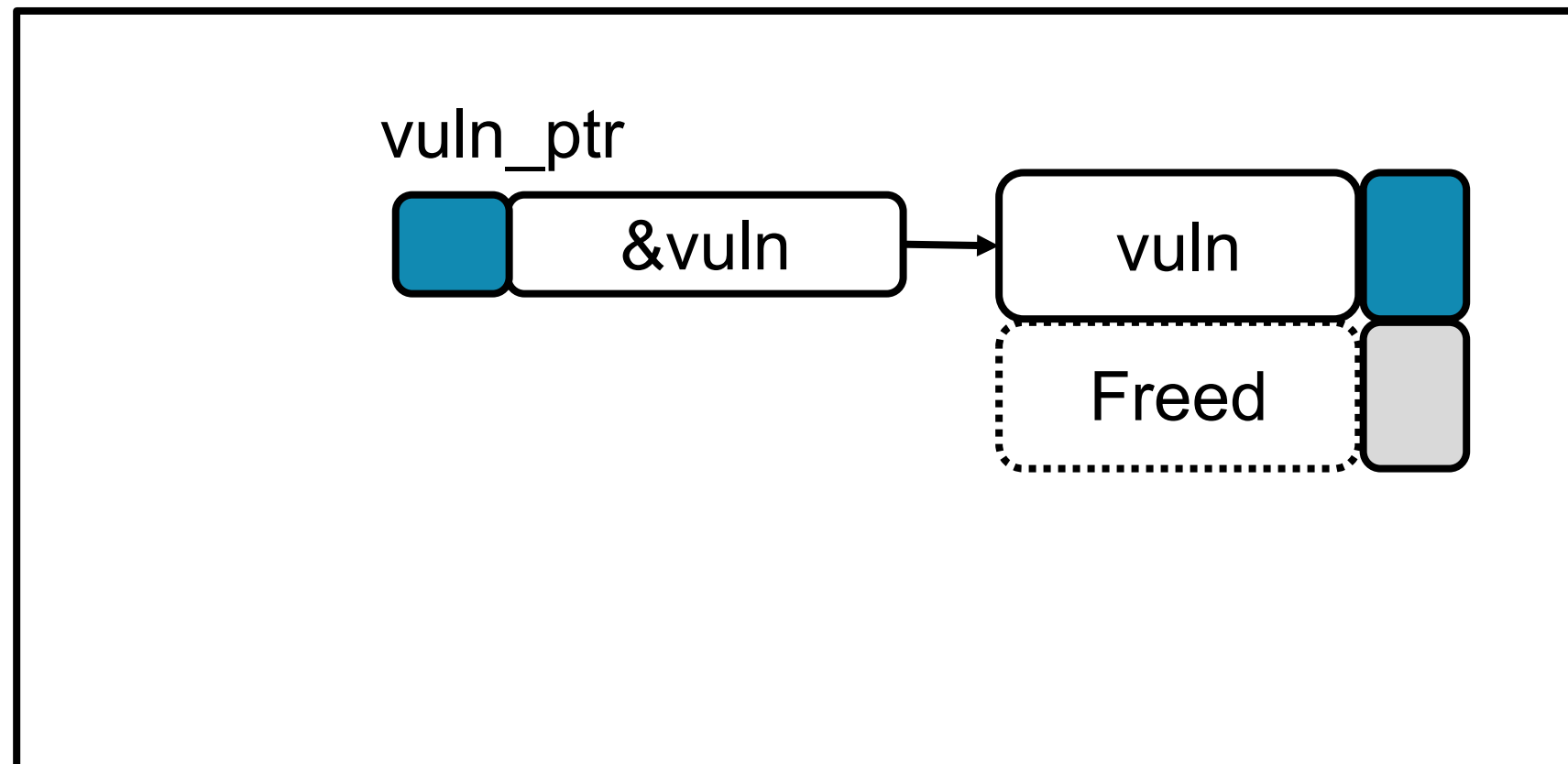


# 3. Reallocate target on tag mismatch




vuln.tag = 

Free(target);

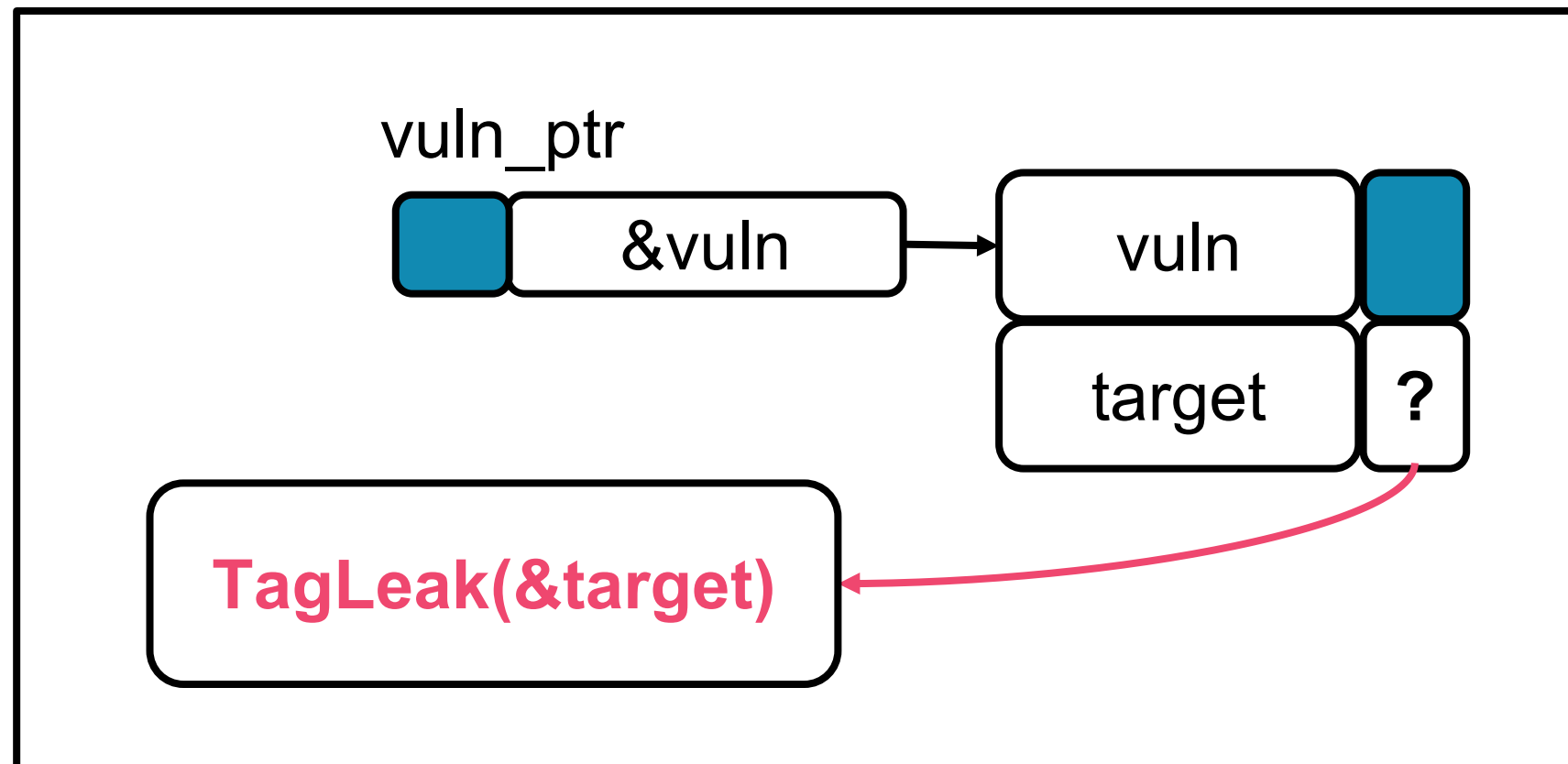


# 3. Reallocate target on tag mismatch



vuln.tag = 



Free(target);  
Alloc(target);

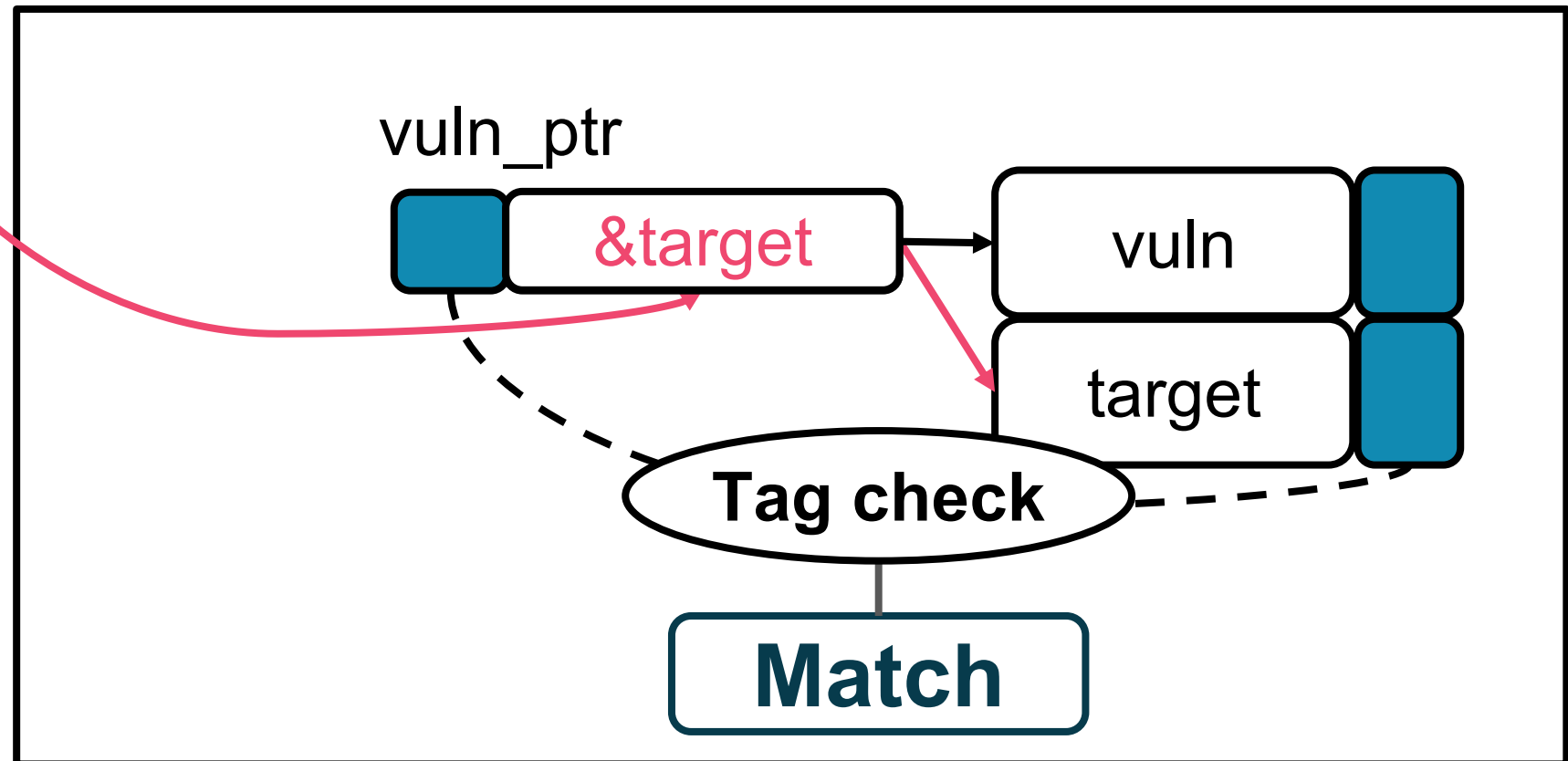


# 4. Trigger vulnerability on tag match

Trigger out-of-bounds access

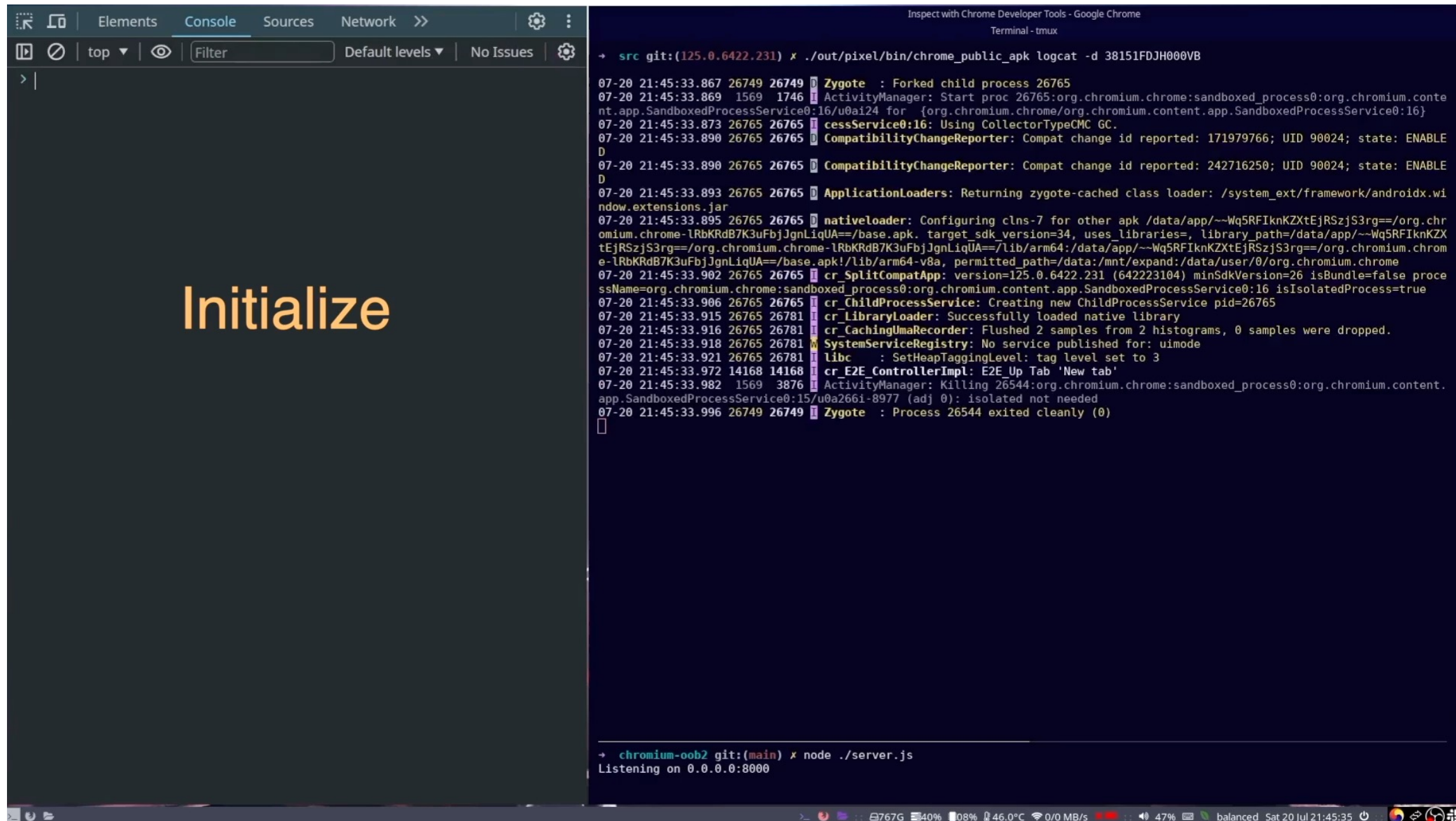


vuln.tag =   
target.tag =   
vuln.tag == target.tag



# CVE-2023-5217 Chrome libvpx heap overflow

## Original Memory Corruption → Attack Fail

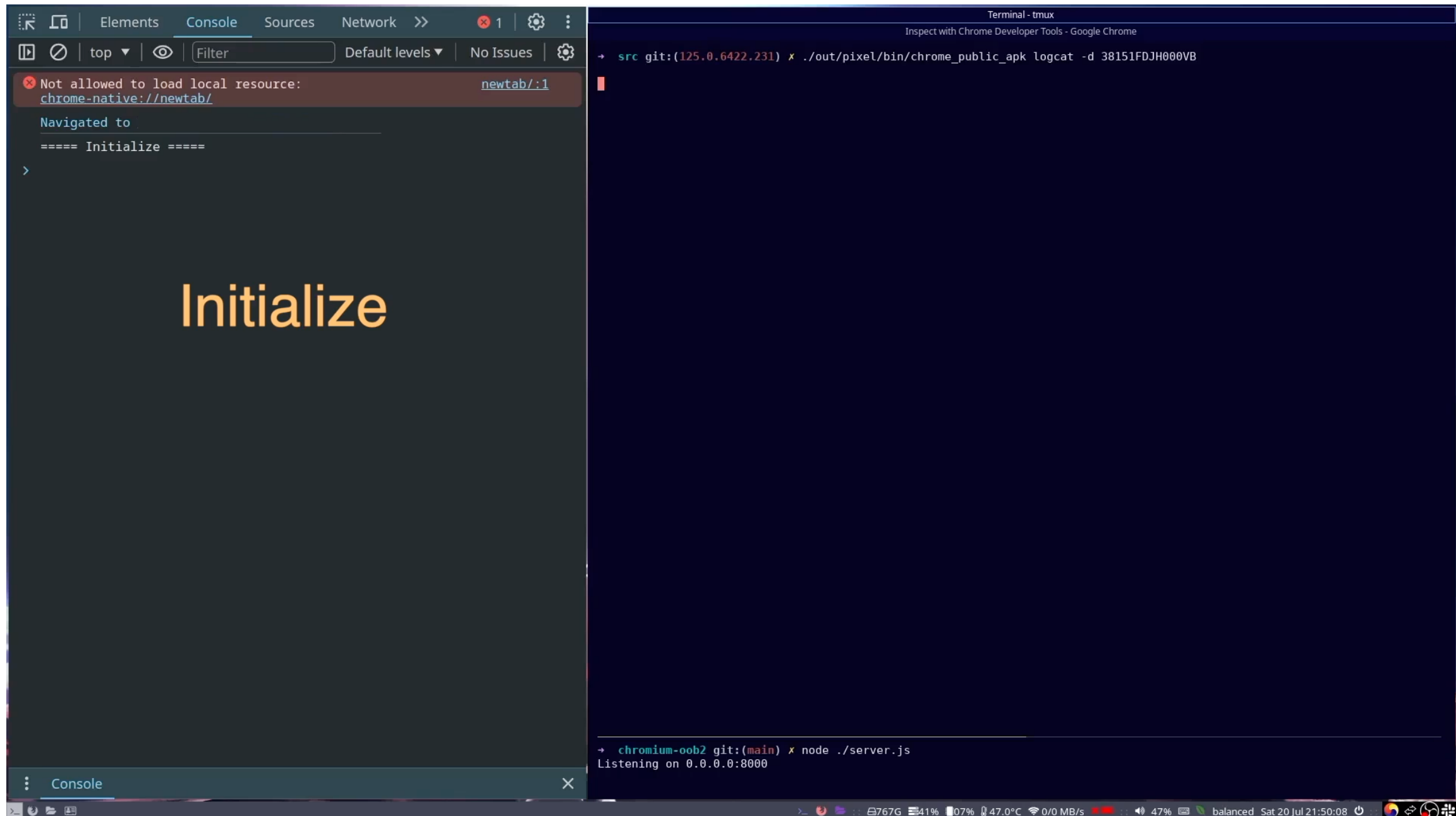


The image shows a screenshot of a Chrome DevTools console and terminal window. The console on the left is titled 'Initialize' and shows a prompt '>'. The terminal on the right shows the execution of a logcat command to monitor a Chrome process. The log output includes various system and application messages, such as 'Zygote : Forked child process 26765', 'ActivityManager: Start proc 26765:org.chromium.chrome:sandboxed\_process0:org.chromium.content.app.SandboxedProcessService0:16/u0ai24 for {org.chromium.chrome/org.chromium.content.app.SandboxedProcessService0:16}', and 'Zygote : Process 26544 exited cleanly (0)'. The terminal prompt at the bottom indicates the process is running in a 'chromium-oob2' directory.

```
src git:(125.0.6422.231) x ./out/pixel/bin/chrome_public_apk logcat -d 38151FDJH000VB
07-20 21:45:33.867 26749 26749 D Zygote : Forked child process 26765
07-20 21:45:33.869 1569 1746 I ActivityManager: Start proc 26765:org.chromium.chrome:sandboxed_process0:org.chromium.conte
nt.app.SandboxedProcessService0:16/u0ai24 for {org.chromium.chrome/org.chromium.content.app.SandboxedProcessService0:16}
07-20 21:45:33.873 26765 26765 I ProcessService0:16: Using CollectorTypeCMC GC.
07-20 21:45:33.890 26765 26765 D CompatibilityChangeReporter: Compat change id reported: 171979766; UID 90024; state: ENABLE
D
07-20 21:45:33.890 26765 26765 D CompatibilityChangeReporter: Compat change id reported: 242716250; UID 90024; state: ENABLE
D
07-20 21:45:33.893 26765 26765 D ApplicationLoaders: Returning zygote-cached class loader: /system_ext/framework/androidx.wi
ndow.extensions.jar
07-20 21:45:33.895 26765 26765 D nativeloader: Configuring clns-7 for other apk /data/app/~Wq5RFIknKZXtEjRSzjS3rg==/org.chr
omium.chrome-lRbKRdB7K3uFbjJgnLiqUA==/base.apk. target_sdk_version=34, uses_libraries=, library_path=/data/app/~Wq5RFIknKZX
tEjRSzjS3rg==/org.chromium.chrome-lRbKRdB7K3uFbjJgnLiqUA==/lib/arm64:/data/app/~Wq5RFIknKZXtEjRSzjS3rg==/org.chromium.chrom
e-lRbKRdB7K3uFbjJgnLiqUA==/base.apk!/lib/arm64-v8a, permitted_path=/data:/mnt/expand:/data/user/0/org.chromium.chrome
07-20 21:45:33.902 26765 26765 I cr_SplitCompatApp: version=125.0.6422.231 (642223104) minSdkVersion=26 isBundle=false proce
ssName=org.chromium.chrome:sandboxed_process0:org.chromium.content.app.SandboxedProcessService0:16 isIsolatedProcess=true
07-20 21:45:33.906 26765 26765 I cr_ChildProcessService: Creating new ChildProcessService pid=26765
07-20 21:45:33.915 26765 26781 I cr_LibraryLoader: Successfully loaded native library
07-20 21:45:33.916 26765 26781 I cr_CachingUmaRecorder: Flushed 2 samples from 2 histograms, 0 samples were dropped.
07-20 21:45:33.918 26765 26781 W SystemServiceRegistry: No service published for: uimode
07-20 21:45:33.921 26765 26781 I libc : SetHeapTaggingLevel: tag level set to 3
07-20 21:45:33.972 14168 14168 I cr_E2E_ControllerImpl: E2E Up Tab 'New tab'
07-20 21:45:33.982 1569 3876 I ActivityManager: Killing 26544:org.chromium.chrome:sandboxed_process0:org.chromium.content.
app.SandboxedProcessService0:15/u0a266i-8977 (adj 0): isolated not needed
07-20 21:45:33.996 26749 26749 I Zygote : Process 26544 exited cleanly (0)
[]

chromium-oob2 git:(main) x node ./server.js
Listening on 0.0.0.0:8000
```

# CVE-2023-5217 Chrome libvpx heap overflow With MTE Tag Leakage → Attack Success



The image shows a screenshot of a web browser's developer tools console and a terminal window. The console on the left displays a message: "Not allowed to load local resource: chrome-native://newtab/" followed by "Navigated to" and "==== Initialize =====". The word "Initialize" is written in large orange text over the console. The terminal window on the right shows a command: "src git:(125.0.6422.231) x ./out/pixel/bin/chrome\_public\_apk logcat -d 38151FDJH000VB". At the bottom of the terminal, another command is visible: "chromium-oob2 git:(main) x node ./server.js Listening on 0.0.0.0:8000". The system tray at the bottom indicates the date and time as "Sat 20 Jul 21:50:08".

# Vendor Responses

## ARM

- Acknowledged the MTE tag side-channel in multiple ARM cores
- MTE Tags are not a secret
  - Tag leakage is not a security vulnerability
- Expected the cost of the hardware fix to be low and recommended the fix.

ARM MTE Security Updates:

<https://developer.arm.com/Arm%20Security%20Center/Arm%20Memory%20Tagging%20Extension>



# Vendor Responses

## Google Android Security Team

- MTE tag leakage are **hardware flaw** of Pixel 8 & Pixel 8 pro
- **Still, MTE is a strong mitigation against limited-shot exploits:**
  - Minimal attack surface (e.g., Messaging app)
  - Physically remote attack (e.g., Bluetooth, NFC, Wi-Fi, ...)
  - Process isolation, IPC attack (e.g., Android, Chrome browser)

# Vendor Responses

## Google Chrome V8 Security Team

- **data confidentiality** (including MTE tag's confidentiality) is out of scope of the V8 Sandbox
- Currently doesn't plan to adopt MTE on renderer due to **potential side-channel issues**

# Takeaway

- **ARM MTE** is a promising security feature to defend against **memory corruption attacks**
- However, current MTE hardware contains **tag leakage side-channel issues**
- MTE-based security can be improved by **software and hardware enhancement** in the future

**Questions?**